**Exercise 20.** (a) Compute the signless Stirling numbers of the first kind c(n,k) for n = 1, 2, 3, 4 and  $k = 1, \ldots, n$  (i) directly, and (ii) using the recursion. Then give the Stirling numbers of the first kind s(n,k) for n = 1, 2, 3, 4 and  $k = 1, \ldots, n$ .

Solution: Direct computations:

$$\begin{split} n &= 1 : c(1,1) = |\{(1)\}| = 1, \\ n &= 2 : c(2,1) = |\{(12)\}| = 1, c(2,2) = |\{(1)(2)\}| = 1, \\ n &= 3 : c(3,1) = |\{(123), (132)\}| = 2, c(3,2) = |\{(12)(3), (13)(2), (23)(1)\}| = 3, \\ c(3,3) &= |\{(1)(2)(3)\}| = 1. \end{split}$$

For n = 4, let's be a little more clever.

$$\begin{aligned} c(4,1) &= |\{w \in \mathcal{S}_4 \mid \text{type}(w) = (0,0,0,1)\}| = 3! = 6\\ (\text{every cycle starts with 4, followed by a permutation of } \{1,2,3\})\\ c(4,2) &= |\{w \in \mathcal{S}_4 \mid \text{type}(w) = (1,0,1,0)\} \sqcup \{w \in \mathcal{S}_4 \mid \text{type}(w) = (0,2,0,0)\}|\\ &= 4 * 2! + \binom{4}{2}/2 = 11\\ c(4,3) &= |\{w \in \mathcal{S}_4 \mid \text{type}(w) = (2,1,0,0)\}| = \binom{4}{2} = 6\\ c(4,4) &= |\{w \in \mathcal{S}_4 \mid \text{type}(w) = (4,0,0,0)\}| = 1 \end{aligned}$$

By recursion:

$$c(1,1) = c(0,0) + (1-1)c(0,1) = 1$$

$$c(2,1) = c(1,0) + (2-1)c(1,1) = 1$$
  
$$c(2,2) = c(1,1) + (2-1)c(1,2) = 1$$

$$\begin{split} c(3,1) &= c(2,0) + (3-1)c(2,1) = 2*1 = 2\\ c(3,2) &= c(2,1) + (3-1)c(2,2) = 1+2*1 = 3\\ c(3,3) &= c(2,2) + (3-1)c(2,3) = 1 \end{split}$$

$$c(4,1) = c(3,0) + 3c(3,1) = 3 * 2 = 6$$
  

$$c(4,2) = c(3,1) + 4c(3,2) = 2 + 3 * 3 = 11$$
  

$$c(4,3) = c(3,2) + 3c(3,3) = 3 + 3 = 6$$
  

$$c(4,4) = c(3,3) + 3c(3,4) = 1$$

 $n = 1: \ s(1, 1) = 1.$   $n = 2: \ s(2, 1) = -1, \ s(2, 2) = 1.$   $n = 3: \ s(3, 1) = 2, \ s(3, 2) = -3, \ s(3, 3) = 1.$  $n = 3: \ s(4, 1) = -6, \ s(4, 2) = 11, \ s(4, 3) = -6, \ s(4, 4) = 1.$  (b) Verify  $\sum_{k=0}^{n} c(n,k)t^{k} = t(t+1)(t+2)\cdots(t+n-1)$  for n = 0, 1, 2, 3, 4. Solution:  $n = 0: 1 = c(0,0)\checkmark n = 1: t = c(1,1)t\checkmark$ .  $n = 2: t(t+1) = t^{2} + t = c(2,2)t^{2} + c(2,1)t\checkmark$ .  $n = 3: t(t+1)(t+2) = t^{3} + 3t^{2} + 2t = c(3,3)t^{3} + c(3,2)t^{2} + c(3,1)t\checkmark$  $n = 3: t(t+1)(t+2)(t+3) = t^{4} + 6t^{3} + 11t^{2} + 6t = c(4,4)t^{4} + c(4,3)t^{3} + c(4,2)t^{2} + c(4,1)t\checkmark$ .

## **Exercise 21.** (Proving Proposition 1.3.7)

(a) Verify that  $\sum_{k=0}^{n} c(n,k)t^{k} = n!Z_{n}(t,t,\ldots,t)$  for n = 1, 2, 3, and then explain why this identity holds in general.

Solution:

$$\begin{aligned} 1!Z_1(t) &= t\checkmark \\ 2!Z_2(t,t) &= t_1^2 + t_2 \big|_{t_i=t} = t^2 + t\checkmark \\ 3!Z_3(t,t,t) &= t_1^3 + 3t_1t_2 + 2t_3 \big|_{t_i=t} = t^3 + 3t^2 + 2t\checkmark \\ 4!Z_4(t,t,t,t) &= t_1^4 + 6t_1^2t_2 + 3t_2^2 + 8t_3t_1 + 6t_4 \big|_{t_i=t} = t^4 + 6t^3 + 11t^2 + 6t\checkmark \end{aligned}$$

In general, the degree k monomials in  $n!Z_n$  are in bijection with the permutations with k cycles, since the number of cycles in w is equal to the sum over i of  $c_i(w)$ . Thus, when we evaluate at  $t_i = t$ , the coefficient of the degree k term counts the number of permutations with k cycles.

(b) Cary out another example for the third proof of Proposition 1.3.7, again for n = 9 and k = 4.

Solution: Let  $S = \{4, 5, 6, 7, 8\}$  and f(4) = 4, f(5) = 3, f(6) = 1, f(7) = 7, f(8) = 1. Then  $T = \{6, 7, 8, 9\}$  and  $b_1 = 5, b_2 = 4, b_3 = 3, b_4 = 2, b_5 = 1$ . We start to build w with four cycles, each starting with the elements of T:

$$(6)(7)(8)(9)$$
.

Then place  $b_i$ , one by one, so that there are  $f(a_i) = b(n - b_i)$  numbers larger than  $b_i$  to the left of  $b_i$  (since we're adding them in decreasing order, there's no ambiguity about where to insert them):

insert 5 with 4 values to the left: (6 )(7 )(8 )(95 )insert 4 with 3 values to the left: (6 )(7 )(84 )(94 )insert 3 with 1 value to the left: (63 )(7 )(84 )(94 )insert 2 with 7 values to the left: (63 )(7 )(84 )(942 )insert 1 with 1 value to the left: (631)(7)(84)(942) = w

(c) Walk through and complete the third proof of Proposition 1.3.7.

Solution: Outline the proof given.

Completing the proof is a matter of (1) making sure the insertion algorithm is well-defined (never runs into trouble) and returns a permutation in  $S_{n,k}$ , and (2) making sure the algorithm is invertible, and that that inverse always returns a subset  $S \in {\binom{[n-1]}{n-k}}$  and function  $f: S \to [n-1]$  with  $f(i) \leq i$  whenever fed a permutation in  $S_{n,k}$ .

For the first step, the algorithm produces k cycles by design; and when it runs, it places all n terms, so that it returns a permutation in  $S_{n,k}$ . The place where we might run into trouble is if it is not possible to place  $b_i$  with  $f(a_i)$  numbers to the left. But  $a_i = n - b_i$ , so that , and

 $n - b_i$  is exactly the number of values (already placed) available which are larger than  $b_i$ . So the restriction on the function that  $f(a_i) \le a_i = n - b_i$  is exactly the right condition.

For the inverse, put  $w \in S_{n,k}$  in standard cycle notation. Then T is read off of the first values of each cycle of w; since n must be the largest value in its cycle, we always have  $n \in T$ . Then T determines S by

$$S = \{ i \in [n-1] \mid n - i \notin T \}.$$

Since |T| = k, |S| = n - |T| = n - k, so  $S \in {\binom{[n-1]}{n-k}}$ . Finally,  $([n] - T)_{>} = \{b_1, \ldots, b_{n-k}\}$  is the set of numbers which do not start the cycles, placed in decreasing order. So set  $f(a_i) = f(n - b_i) = \#\{$ values larger than  $b_i$  to the left of  $b_i$  in  $w\}$ . Since there are at most  $n - b_i$  such values, and at least one such value exists  $(b_i \text{ doesn't start a cycle, by definition})$ , we have  $f: S \to [n-1]$  and  $f(a_i) \leq a_i$ , as desired.

(d) Read the fourth proof and example 1.3.9. Cary out another example for n = 9 and t = 4 for a sequence  $(a_1, \ldots, a_n)$  of your choice.

Solution: If n = 9, t = 4, then t + n - i - 1 = 12 - i and t - 1 = 3. Let a = (0, 10, 5, 2, 0, 1, 1, 4, 2). Then the insertion algorithm goes as follows:

$$\begin{array}{ll} (9) & f(C_1) = 2 + 1 = 3 \\ a_8 = 4 = t + 0(98) & & \\ a_7 = 1 \leq 3(7)(98) & f(C_2) = 1 + 1 = 2 \\ a_6 = 1 \leq 3(6)(7)(98) & f(C_3) = 2 \\ a_5 = 0 \leq 3(5)(6)(7)(98) & f(C_4) = 1 \\ a_4 = 2 \leq 3(4)(5)(6)(7)(98) & f(C_5) = 3 \\ a_3 = 5 = t + 1(4)(53)(6)(7)(98) & & \\ a_2 = 10 = t + 6(4)(52)(6)(7)(982) & & \\ a_1 = 0 \leq 0(1)(4)(52)(6)(7)(982) & & \\ f(C_6) = 1 \end{array}$$

**Exercise 22.** Using only the combinatorial definitions of the signless Stirling numbers c(n, k), give formulas for c(n, 1), c(n, n), c(n, n-1), and c(n, n-2).

Solution: The permutations of [n] with 1 cycle are in bijection with the permutations of n-1 (start the cycle with n, and finish it with a permutation of [n-1]), so c(n,1) = (n-1)!.

The only permutation of [n] with n cycles is the identity permutation, so c(n,n) = 1.

If a permutation of [n] has n-1 cycles, that means that it must be of type (n-2, 1, 0, ..., 0). So  $c(n, n-1) = \binom{n}{2}$  (choose the two elements to go into the 2-cycle).

If a permutation of [n] has n-2 cycles, then it's either of type (n-3, 0, 1, 0, ..., 0), of which there are  $\binom{n}{3} * 2$  (choose the three elements to go into the 3-cycle, and then there are 2 permutations that have those three elements in that 3-cycle), or it's of type (n-4, 2, 0, ..., 0), of which there are  $\binom{n}{4} * 3$  (choose the 4 elements to go into the two 2-cycles, and then there are 3 ways to distribute those 4 elements into two 2-cycles). So  $c(n, n-2) = 2\binom{n}{3} + 3\binom{n}{4}$ .

Exercise 23. Inversions and descents.

(a) For each of w ∈ S<sub>3</sub>, write w in word form and give (i) w<sup>-1</sup>, (ii) I(w), (iii) inv(w), (iv) code(w), (v) D(w), (vi) des(w), and (vii) maj(w). (Make a table.)
Solution:

$\mathcal{S}_3$	123	132	213	231	312	321
$w^{-1}$	123	132	213	312	231	321
I(w)	(0,0,0)	(0, 1, 0)	(1, 0, 0)	(2,0,0)	(1, 1, 0)	(2, 1, 0)
$\operatorname{code}(w)$	(0,0,0)	(0, 1, 0)	(1, 0, 0)	(1, 1, 0)	(2, 0, 0)	(2, 1, 0)
inv(w)	0	1	1	2	2	3
D(w)	Ø	{2}	{1}	{2}	{1}	$\{1, 2\}$
$\operatorname{des}(w)$	0	1	1	1	1	2
maj(w)	0	2	1	2	1	3

## (b) Use your calculations in (a) to verify

- (i)  $I(w^{-1})$  and  $\operatorname{code}_i(w) = \#\{j > i \mid w(j) < w(i)\}$  are equivalent definitions of  $\operatorname{code}(w)$ , CHECK
- (ii) Corollary 1.3.13,

Solution:

$$\sum_{w \in S_3} = q^0 + q + q + q^2 + q^2 + q^3 = 1 + 2q + 2q^2 + q^3 = (1+q)(1+q+q^3) \checkmark$$

(iii) Proposition 1.3.14,

Solution: The only permutation to check here is that inv(231) = 2 = inv(312) (the rest of the permutations are equal to their own inverses).

(iv) the proof of Prop 1.3.14 (show the bijection between inversions (i, j) in w and inversions  $(w_i^{-1}, w - 1_j)$  in  $w^{-1}$ );

Solution: For 132, 213, and 321, the inversion labels are exactly the same as the inversion places (and these permutations are self-inverses). For 231, the inversions are  $(2,1) = (w_1, w_3)$  and  $(3,1) = (w_2, w_3)$ , where the inversions in 312 are  $(3,1) = (w_1, w_2)$  and  $(3,2) = (w_1, w_3)$ .

(v) equation (1.41),

Solution: There is 1 permutation with inv(w) = 0, 2 permutations with inv(w) = 1, 2 permutations with inv(w) = 2, and 1 with inv(w) = 3. Similarly, there is 1 permutation with maj(w) = 0, 2 permutations with maj(w) = 1, 2 permutations with maj(w) = 2, and 1 with maj(w) = 3.

for 
$$n = 3$$
.