**Exercise 32.** Complete the proof of the following theorem: The following are equivalent for a simple graph G.

- (1) G is a tree.
- (2) G is a minimal connected graph, i.e. every edge in E is a cut edge.
- (3) G is a maximally acyclic graph, i.e. G is acyclic, and adding any edge between two nonadjacent vertices creates a cycle.

Solution: (1)  $\implies$  (2): Let  $G = (V, E, \phi)$  be a tree, so that it is connected and acyclic. Let  $e \in E$ , and let  $\phi(e) = \{u, v\}$ . Suppose there is some path P from u to v that does not go through e. Then extending P by e results in a cycle, which is a contradiction. Thus e is a cut edge, and so G is minimally connected.

(2)  $\implies$  (1): Let G be a minimally connected graph. Suppose G contains a cycle C, with e an edge in C. Then any walk which goes through e can be rerouted along C - e. Thus e is not a cut edge, a contradiction. Thus G is connected and acyclic, and therefore a tree.

(1)  $\implies$  (3): Let  $G = (V, E, \phi)$  be a tree, so that it is connected and acyclic. Consider an edge  $e \in \overline{G}$ , with  $\phi(e) = \{u, v\}$ . Since G is connected, there is a path P in G from u to v. Thus, in G + e, extending P by e results in a cycle. Thus G is maximally acyclic.

Not (1)  $\implies$  not (3): Let  $G = (V, E, \phi)$  be a graph that is not a tree. Then either G contains at least one cycle, implying G is not acyclic, or G is a forrest with at least two connected components. If G is a forrest, then let u and v be in different connected components. Thus  $uv \notin E$ , and there are no paths from u to v. Thus G + e is also acyclic. In either case, G is not maximally acyclic.

**Exercise 33.** (a) How many spanning trees does  $C_5$  have?

Solution: One corresponding to the deletion of each edge in  $C_5$ : 5.

(b) Let



(i) Calculate diam(G) and rad(G).

Solution:

d(u,v)	a	b	c	d	e	f	g	h
a	0	2	2	1	1	3	1	3
b	2	0	2	3	1	1	3	3
С	2	2	0	3	1	1	1	1
d	1	3	3	0	2	4	2	4
e	1	1	1	2	0	2	2	2
f	3	1	1	4	2	0	2	2
g	1	3	1	2	2	2	0	2
h	3	3	1	4	2	2	2	0
		·	·	·			·	

 $diam(G) = d(d, h) = 4, \qquad rad(G) = 2 \text{ (see } e)$ 

(ii) Fix the vertex a, and give  $V_i = \{u \in V \mid d(u, a) = i\}$ . Solution:

```
V_0 = \{a\}

V_1 = \{d, e, g\}

V_2 = \{b, c\}

V_3 = \{f, h\}.
```

- (iii) Build a spanning tree using Method 1 from the notes using v = a (show your steps!). What is the radius of the resulting tree?'
- (iv) Find a central vertex v, i.e. one for which  $\max_{u \in V} d(u, v) = \operatorname{rad}(G)$ , and build a spanning tree using Method 1 from the notes using that vertex (show your steps!). What is the radius of the resulting tree?
- (v) Build a spanning tree using Method 2 from the notes starting with  $T_1$  being the isolated vertex *a* (show your steps!).