## Solutions for HW10

Exercise 32. Complete the proof of the following theorem:
The following are equivalent for a simple graph $G$.
(1) $G$ is a tree.
(2) $G$ is a minimal connected graph, i.e. every edge in $E$ is a cut edge.
(3) $G$ is a maximally acyclic graph, i.e. $G$ is acyclic, and adding any edge between two nonadjacent vertices creates a cycle.

Solution: $(1) \Longrightarrow(2)$ : Let $G=(V, E, \phi)$ be a tree, so that it is connected and acyclic. Let $e \in E$, and let $\phi(e)=\{u, v\}$. Suppose there is some path $P$ from $u$ to $v$ that does not go through $e$. Then extending $P$ by $e$ results in a cycle, which is a contradiction. Thus $e$ is a cut edge, and so $G$ is minimally connected.
$(2) \Longrightarrow(1)$ : Let $G$ be a minimally connected graph. Suppose $G$ contains a cycle $C$, with $e$ an edge in $C$. Then any walk which goes through $e$ can be rerouted along $C-e$. Thus $e$ is not a cut edge, a contradiction. Thus $G$ is connected and acyclic, and therefore a tree.
$(1) \Longrightarrow(3)$ : Let $G=(V, E, \phi)$ be a tree, so that it is connected and acyclic. Consider an edge $e \in \bar{G}$, with $\phi(e)=\{u, v\}$. Since $G$ is connected, there is a path $P$ in $G$ from $u$ to $v$. Thus, in $G+e$, extending $P$ by $e$ results in a cycle. Thus $G$ is maximally acyclic.
Not (1) $\Longrightarrow$ not (3): Let $G=(V, E, \phi)$ be a graph that is not a tree. Then either $G$ contains at least one cycle, implying $G$ is not acyclic, or $G$ is a forrest with at least two connected components. If $G$ is a forrest, then let $u$ and $v$ be in different connected components. Thus $u v \notin E$, and there are no paths from $u$ to $v$. Thus $G+e$ is also acyclic. In either case, $G$ is not maximally acyclic.

Exercise 33. (a) How many spanning trees does $C_{5}$ have?
Solution: One corresponding to the deletion of each edge in $C_{5}: 5$.
(b) Let

(i) Calculate $\operatorname{diam}(G)$ and $\operatorname{rad}(G)$.

Solution:

| $d(u, v)$ | $a$ | $b$ | $c$ | $d$ | $e$ | $f$ | $g$ | $h$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | 0 | 2 | 2 | 1 | 1 | 3 | 1 | 3 |
| $b$ | 2 | 0 | 2 | 3 | 1 | 1 | 3 | 3 |
| $c$ | 2 | 2 | 0 | 3 | 1 | 1 | 1 | 1 |
| $d$ | 1 | 3 | 3 | 0 | 2 | 4 | 2 | 4 |
| $e$ | 1 | 1 | 1 | 2 | 0 | 2 | 2 | 2 |
| $f$ | 3 | 1 | 1 | 4 | 2 | 0 | 2 | 2 |
| $g$ | 1 | 3 | 1 | 2 | 2 | 2 | 0 | 2 |
| $h$ | 3 | 3 | 1 | 4 | 2 | 2 | 2 | 0 |

$$
\operatorname{diam}(G)=d(d, h)=4, \quad \operatorname{rad}(G)=2(\text { see } e)
$$

(ii) Fix the vertex $a$, and give $V_{i}=\{u \in V \mid d(u, a)=i\}$.

Solution:

$$
\begin{aligned}
V_{0} & =\{a\} \\
V_{1} & =\{d, e, g\} \\
V_{2} & =\{b, c\} \\
V_{3} & =\{f, h\} .
\end{aligned}
$$

(iii) Build a spanning tree using Method 1 from the notes using $v=a$ (show your steps!). What is the radius of the resulting tree?'
(iv) Find a central vertex $v$, i.e. one for which $\max _{u \in V} d(u, v)=\operatorname{rad}(G)$, and build a spanning tree using Method 1 from the notes using that vertex (show your steps!). What is the radius of the resulting tree?
(v) Build a spanning tree using Method 2 from the notes starting with $T_{1}$ being the isolated vertex $a$ (show your steps!).

