Lecture 4: Divisibility and Greatest Common Divisor

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 since $6 = 3 * 2;$
 $15|60$ since $60 = 15 * 4;$

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 $\begin{array}{rll} 3|6 & \text{since} & 6=3*2;\\ 15|60 & \text{since} & 60=15*4;\\ 15 \nmid 25 & \text{since there is no } m \in \mathbb{Z} \text{ such that } 25=15*m. \end{array}$

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 In general, for any $n \in \mathbb{Z}$,

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Example: the divisors of 12 are 1, 2, 3, 4, 6, and 12.

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For example, 3 is a divisor of 30, but not 40; 4 is a divisor of 40, but not 30; 1, 2, 5, and 10 are all common divisors of 30 and 40. The greatest common divisor of a and b, denoted gcd(a, b) is largest number that divides both a and b.

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Always, gcd(a, b) = gcd(b, a). If b|a, then gcd(a, b) = b. If gcd(a, b) = 1, we say that a and b are relatively prime. Example:

The divisors of 25 are 1, 5, and 25; the divisors of 12 are 1, 2, 3, 4, 6, and 12; so 25 and 12 are relatively prime (even though neither is prime).

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$$\begin{array}{rl} 19500=2^2*3*5^3*13 \quad \text{and} \quad 440=2^3*5*11,\\ \text{so} \quad \gcd(19500,400)=2^2*5=\fbox{20}. \end{array}$$

In other words, gcd(a, b) will be the product over primes p to the highest power n such that $p^n|a$ and $p^n|b$.

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You try: compute the prime factorizations of $12, 30, 35, {\rm and}\ 84, {\rm and}\ use them to compute$

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Not computationally efficient either! (Prime factorization is computationally difficult/not possible without a list of primes.)

First, we'll need the division algorithm, which says for any $a, b \in \mathbb{Z}$ with $b \neq 0$, there are unique integers q and r satisfying

$$a = bq + r$$
 and $0 \leq r < |b|$.

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Proof: (sketch) If a and b are the same sign, subtract b from a until the result is between 0 and |b| - 1. The result is r and the number of subtractions is q. If they're different signs, add b to a until the result is between 0 and |b| - 1. The result is r and the number of additions is -q.

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If $a_2 = 5, b_2 = 2$, then $q_2 = 2$ and $r_2 = 1$ since $5 = 2 * 2 + 1$.

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Play this game again with new a and b:

- 1. Start with $a_1 = a$ and $b_1 = b$.
- 2. Compute the quotient q_i and remainder r_i in dividing a_i by b_i .
- 3. Repeat the division algorithm using $a_i = b_{i-1}$ and $b_i = r_{i-1}$.
- 4. Iterate until you get $r_n = 0$. Then compare gcd(a, b) with r_{n-1} .

You try: Do this process with a = 30, b = 12, and then with a = 84, b = 30.

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You try: Do this process with a = 30, b = 12, and then with a = 84, b = 30. Claim: If n is the first time that $r_n = 0$, then $r_{n-1} = \text{gcd}(a, b)$. We have

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- 4. Iterate until you get $r_n = 0$. Then compare gcd(a, b) with r_{n-1} .

You try: Do this process with a = 30, b = 12, and then with a = 84, b = 30. Claim: If n is the first time that $r_n = 0$, then $r_{n-1} = \gcd(a, b)$. Note that if r = 0 in the first step, then b|n, so $\gcd(a, b) = b$.

Spreadsheet functions

For a and integer and b a positive integer, =FLOOR(a, b)gives the largest multiple of b less or equal to a. Namely, if a = bq + r, then FLOOR(a, b) = bq.

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=FLOOR(17,5) returns 15, =FLOOR(-17,5) returns -20, =FLOOR(17,-5) returns an error.

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gives the largest multiple of b less or equal to a. Namely, if a=bq+r, then $\mathsf{FLOOR}(a,b)=bq.$

Example:

= FLOOR(17, 5) returns 15, = FLOOR(-17, 5) returns -20, = FLOOR(17, -5) returns an error.So to compute q and r such that a = bq + r, = FLOOR(a, b)/b returns q,= a - FLOOR(a, b) returns r.

In general, our process looks like

$$a = b * q_{1} + r_{1}$$

$$b = r_{1} * q_{2} + r_{2}$$

$$r_{1} = r_{2} * q_{3} + r_{3}$$

$$\vdots$$

$$r_{n-4} = r_{n-3} * q_{n-2} + r_{n-2}$$

$$r_{n-3} = r_{n-2} * q_{n-1} + r_{n-1} \leftarrow \gcd(a, b)?$$

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To make everything look the same, let $r_{-1} = a$ and $r_0 = b$. So every line comes in the form

$$r_{i-2} = r_{i-1} * q_i + r_i.$$

Let $r_{-1} = a$ and $r_0 = b$, so that the algorithm looks like

$$\begin{array}{rcrcrcrcrc} r_{-1} & = & r_0 * q_1 & + & r_1 \\ r_0 & = & r_1 * q_2 & + & r_2 \\ r_1 & = & r_2 * q_3 & + & r_3 \\ & \vdots & & & \\ r_{n-4} & = & r_{n-3} * q_{n-2} & + & r_{n-2} \\ r_{n-3} & = & r_{n-2} * q_{n-1} & + & r_{n-1} & \leftarrow \gcd(a, b)? \\ r_{n-2} & = & r_{n-1} * q_n & + & 0 & \leftarrow r_n \end{array}$$

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Let $r_{-1} = a$ and $r_0 = b$, so that the algorithm looks like

$$\begin{array}{rcrcrcrcrc} r_{-1} &=& r_{0} * q_{1} &+& r_{1} \\ r_{0} &=& r_{1} * q_{2} &+& r_{2} \\ r_{1} &=& r_{2} * q_{3} &+& r_{3} \\ &\vdots \\ r_{n-4} &=& r_{n-3} * q_{n-2} &+& r_{n-2} \\ r_{n-3} &=& r_{n-2} * q_{n-1} &+& r_{n-1} &\leftarrow \gcd(a,b)? \\ r_{n-2} &=& r_{n-1} * q_{n} &+& 0 &\leftarrow r_{n} \end{array}$$

Last line: $r_{n-2} = r_{n-1}q_n$. So $r_{n-3} = r_{n-2}q_{n-1} + r_{n-1} = (r_{n-1}q_n)q_{n-1} + r_{n-1} = r_{n-1}(q_nq_{n-1} + 1)$. Then

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Then

 $r_{n-4} = r_{n-3}q_{n-2} + r_{n-2} = r_{n-1}(q_nq_{n-1} + 1)q_{n-2} + r_{n-1}q_n$

Let $r_{-1} = a$ and $r_0 = b$, so that the algorithm looks like

$$\begin{array}{rcrcrcrcrc} r_{-1} &=& r_{0} * q_{1} &+& r_{1} \\ r_{0} &=& r_{1} * q_{2} &+& r_{2} \\ r_{1} &=& r_{2} * q_{3} &+& r_{3} \\ &\vdots \\ r_{n-4} &=& r_{n-3} * q_{n-2} &+& r_{n-2} \\ r_{n-3} &=& r_{n-2} * q_{n-1} &+& r_{n-1} &\leftarrow \gcd(a,b)? \\ r_{n-2} &=& r_{n-1} * q_{n} &+& 0 &\leftarrow r_{n} \end{array}$$

Last line: $r_{n-2} = r_{n-1}q_n$. So $r_{n-3} = r_{n-2}q_{n-1} + r_{n-1} = (r_{n-1}q_n)q_{n-1} + r_{n-1} = r_{n-1}(q_nq_{n-1} + 1)$. Then

$$r_{n-4} = r_{n-3}q_{n-2} + r_{n-2} = r_{n-1}(q_nq_{n-1}+1)q_{n-2} + r_{n-1}q_n$$

= $r_{n-1}(q_nq_{n-1}q_{n-2} + q_{n-2} + 1).$

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.
So
 $r_{n-3} = r_{n-2}q_{n-1} + r_{n-1} = (r_{n-1}q_n)q_{n-1} + r_{n-1} = r_{n-1}(q_nq_{n-1} + 1)$.
Then

$$\begin{aligned} r_{n-4} &= r_{n-3}q_{n-2} + r_{n-2} = r_{n-1}(q_nq_{n-1}+1)q_{n-2} + r_{n-1}q_n \\ &= r_{n-1}(q_nq_{n-1}q_{n-2}+q_{n-2}+1). \end{aligned} \text{And so on} \ldots$$

Example: We saw

$$84 = 30 * 2 + 24$$
$$30 = 24 * 1 + 6$$
$$24 = 6 * 4 + 0.$$

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$$30 = 24 * 1 + 6 = (6 * 4) * 1 + 6 = 6(4 * 1 + 1) = 6 * 5$$

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84 = 30 * 2 + 24 30 = 24 * 1 + 6 24 = 6 * 4 + 0. $r_{n-1} = 6$

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So 6 is a common divisor of 84 and 30.

100 = 36 * 2 + 28 36 = 28 * 1 + 8 28 = 8 * 3 + 48 = 4 * 2 + 0.

$$100 = 36 * 2 + 28$$

$$36 = 28 * 1 + 8$$

$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

$$100 = 36 * 2 + 28$$

$$36 = 28 * 1 + 8$$

$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

So

28 = 8 * 3 + 4

$$100 = 36 * 2 + 28$$

$$36 = 28 * 1 + 8$$

$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

$$28 = 8 * 3 + 4 = (4 * 2) * 3 + 4$$

$$100 = 36 * 2 + 28$$

$$36 = 28 * 1 + 8$$

$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

$$28 = 8 * 3 + 4 = (4 * 2) * 3 + 4 = 4(2 * 3 + 1) = 4 * 7$$

$$100 = 36 * 2 + 28$$

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$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

$$28 = 8 * 3 + 4 = (4 * 2) * 3 + 4 = 4(2 * 3 + 1) = 4 * 7$$

$$36 = 28 * 1 + 8$$

$$100 = 36 * 2 + 28$$

$$36 = 28 * 1 + 8$$

$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

$$28 = 8 * 3 + 4 = (4 * 2) * 3 + 4 = 4(2 * 3 + 1) = 4 * 7$$

$$36 = 28 * 1 + 8 = (4 * 7) * 1 + (4 * 2)$$

$$100 = 36 * 2 + 28$$

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$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

$$28 = 8 * 3 + 4 = (4 * 2) * 3 + 4 = 4(2 * 3 + 1) = 4 * 7$$

$$36 = 28 * 1 + 8 = (4 * 7) * 1 + (4 * 2) = 4(7 * 1 + 2) = 4 * 9$$

$$100 = 36 * 2 + 28$$

$$36 = 28 * 1 + 8$$

$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

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So

$$28 = 8 * 3 + 4 = (4 * 2) * 3 + 4 = 4(2 * 3 + 1) = 4 * 7$$

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$$r_{n-1} = 4$$

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$$28 = 8 * 3 + 4 = (4 * 2) * 3 + 4 = 4(2 * 3 + 1) = 4 * 7$$

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$$100 = 36 * 2 + 28 = (4 * 9) * 2 + (4 * 7) = 4(9 * 2 + 7) = 4 * 25.$$

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$$r_{n-1} = 4$$

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So 4 is a common divisor of 100 and 36.

$$100 = 36 * 2 + 28$$

$$36 = 28 * 1 + 8$$

$$28 = 8 * 3 + 4$$

$$8 = 4 * 2 + 0.$$

$$r_{n-1} = 4$$

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So 4 is a common divisor of 100 and 36.

You try: use the following computations, working backwards, to show that 2 is a common divisor of 100 and 26: 100 = 26 * 3 + 22 26 = 22 * 1 + 4

$$100 = 26 * 3 + 22 22 = 4 * 5 + 2 26 = 22 * 1 + 4 4 = 2 * 2 + 0$$

Letting $r_{-1} = a$ and $r_0 = b$, and computing

we can reverse this process to show that r_{n-1} is, at the very least, a *common divisor* to $a = r_{-1}$ and $b = r_0$.

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$$b = r_0 > r_1 > r_2 > \dots \ge 0.$$

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So since the r_i 's are all *integers*, this process ends at some point.

We have that r_{n-1} is a common divisor to a an b. Now why is it the *greatest* common divisor?

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$$a = b * q_1 + r_1$$

$$b = r_1 * q_2 + r_2$$

$$r_1 = r_2 * q_3 + r_3$$

$$\vdots$$

$$r_{n-3} = r_{n-2} * q_{n-1} + r_{n-1}$$

$$r_{n-2} = r_{n-1} * q_n + 0$$

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dlpha	=	$deta * q_1$	+	r_1
b	=	$r_1 * q_2$	+	r_2
r_1	=	$r_2 * q_3$	+	r_3
	:			
r_{n-3}	=	$r_{n-2} * q_{n-1}$	+	r_{n-1}
r_{n-2}	=	$r_{n-1} * q_n$	+	0

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Back to our division calculation, and substitute these equations in:

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So d is a divisor of r_{n-1} .

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The Euclidean algorithm for computing the greatest common divisor of two positive numbers a and b is the process or successively dividing until just before you get a 0 divisor (like we just did). Namely, we have the following theorem.

Theorem (Euclidean algorithm)

To compute the greatest common divisor of two positive integers a and b, let $r_{-1} = a$ and $r_0 = b$, and compute successive quotients and remainders

$$r_{i-2} = r_{i-1}q_i + r_i$$

for i = 1, 2, 3, ..., until some remainder r_n is 0. The last nonzero remainder r_{n-1} is then the greatest common divisor of a and b.

This takes at most b steps (actually less), and is *much* more computationally efficient than the other methods.