

Public key cryptography

For **Person A** to receive messages. . .

Step 1: Make a cypher, turning letters into numbers.

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>	<i>G</i>	<i>H</i>	<i>I</i>	<i>J</i>	<i>K</i>	<i>L</i>	<i>M</i>
11	12	13	14	15	16	17	18	19	20	21	22	23

<i>N</i>	<i>O</i>	<i>P</i>	<i>Q</i>	<i>R</i>	<i>S</i>	<i>T</i>	<i>U</i>	<i>V</i>	<i>W</i>	<i>X</i>	<i>Y</i>	<i>Z</i>
24	25	26	27	28	29	30	31	32	33	34	35	36

For example, "TO BE OR NOT TO BE" becomes

30251215252824253030251215:

<i>T</i>	<i>O</i>	<i>B</i>	<i>E</i>	<i>O</i>	<i>R</i>	<i>N</i>	<i>O</i>	<i>T</i>	<i>T</i>	<i>O</i>	<i>B</i>	<i>E</i>
30	25	12	15	25	28	24	25	30	30	25	12	15

(Use a larger key to include spaces and punctuation, etc.)

For **Person A** to receive messages. . .

Step 1: Make a cypher, turning letters into numbers.

Step 2: Choose two large primes p and q , and let $n = pq$.

[Generally, chosen randomly, but within a couple of digits in length of each other.]

Ex: Let $p = 12553$, $q = 13007$. So $n = 163276871$.

Step 3: Compute $\phi(pq) = (p - 1)(q - 1)$, and pick a number k relatively prime to $\phi(pq)$.

[Specifically, pick k randomly between 10 and $\text{lcm}(p - 1, q - 1)$.]

Ex: $\phi(n) = 12552 * 13006 = 163251312$. Pick $k = 79921$.

Step 4: Publish the cypher, n , and k publicly; keep p , q , and $\phi(pq)$ secret.

Published: $n = 163276871$, $k = 79921$, and cypher

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>	<i>G</i>	<i>H</i>	<i>I</i>	<i>J</i>	<i>K</i>	<i>L</i>	<i>M</i>
11	12	13	14	15	16	17	18	19	20	21	22	23
<i>N</i>	<i>O</i>	<i>P</i>	<i>Q</i>	<i>R</i>	<i>S</i>	<i>T</i>	<i>U</i>	<i>V</i>	<i>W</i>	<i>X</i>	<i>Y</i>	<i>Z</i>
24	25	26	27	28	29	30	31	32	33	34	35	36

For **Person B** to send messages. . .

Step 1: Using the provided cypher, turn letters into numbers, and break into small-ish pieces (fewer digits than n).

Ex: TO BE OR NOT TO BE \rightarrow 30251215252824253030251215

$$a_1 = 30251215, \quad a_2 = 25282425, \quad a_3 = 30302512, \quad a_4 = 15.$$

Step 2: Compute $a_i^k \pmod{n}$ for each piece.

Ex:

$$30251215^{79921} \equiv_{163276871} 149419241$$

$$25282425^{79921} \equiv_{163276871} 62721998$$

$$30302512^{79921} \equiv_{163276871} 118084566$$

$$15^{79921} \equiv_{163276871} 40481382$$

Step 3: Send the results.

Back to **Person A:** You know $n = 163276871$, $k = 79921$, and the cypher. You *also* know $p = 12553$ and $q = 13007$, so that $\phi(n) = 163251312$. So now, given that

$$a_1^k \equiv_n 149419241, \quad a_2^k \equiv_n 62721998, \\ a_3^k \equiv_n 118084566, \quad \text{and} \quad a_4^k \equiv_n 40481382,$$

you can use the methods from last time to solve for a_1, a_2, a_3 , and a_4 ! Namely, you use the Euclidean algorithm to compute

$$1 = 145604785 \cdot k - 71282 \cdot \phi(n).$$

Then, you know if $a^k \equiv_n b$, then $a \equiv_n b^u$, where $u = 145604785$.

So, using the method of successive squaring, you are able to compute

$$a_1 \equiv 149419241^{145604785} \equiv 30251215 \pmod{163276871},$$

$$a_2 \equiv 62721998^{145604785} \equiv 25282425 \pmod{163276871},$$

$$a_3 \equiv 118084566^{145604785} \equiv 30302512 \pmod{163276871},$$

$$a_4 \equiv 40481382^{145604785} \equiv 15 \pmod{163276871},$$

as desired.

RSA public key cryptosystem, after Ron Rivest, Adi Shamir, and Leonard Adleman.