

Notation: for a fixed n , let \bar{a} be the least residue of $a \pmod{n}$, i.e. the unique number between 0 and $n - 1$ congruent to a .

Last time: Method of successive squaring

Given x, k , and big n , compute $x^k \pmod{n}$ as follows.

1. If $\gcd(x, n) = 1$, first reduce $k \equiv \bar{k} \pmod{\phi(n)}$, so that by Euler's formula

$$x^k \equiv x^{\bar{k}} \pmod{n}.$$

$$\text{(since } x^k = x^{m\phi(n) + \bar{k}} = (x^{\phi(n)})^m x^{\bar{k}} \equiv_n 1^m \cdot x^{\bar{k}} \text{)}$$

2. Decompose \bar{k} (or k if $\gcd(x, n) \neq 1$) into powers of 2:

$$\bar{k} = 2^{a_1} + 2^{a_2} + \dots + 2^{a_\ell}.$$

3. Use successive squaring (*square, reduce, square, reduce, ...*) to compile a table of data for $x^{2^a} \pmod{n}$, for as many a as you need.

$$(x = \boxed{x^1}, (x^1)^2 = \boxed{x^2}, (x^2)^2 = \boxed{x^4}, (x^4)^2 = \boxed{x^8}, (x^8)^2 = \boxed{x^{16}} \dots)$$

4. Use your table and your decomposition to compute $x^k \pmod{n}$:

$$x^k \equiv \overline{x^{2^{a_1}}} \cdot \overline{x^{2^{a_2}}} \dots \overline{x^{2^{a_\ell}}} \pmod{n}.$$

Goal: Reverse the process. Namely, given k, b , and big n , solve $x^k \equiv b \pmod{n}$ for x .

Process: Assume $\gcd(b, n) = 1 = \gcd(k, \phi(n))$.

1. Compute $\phi(n)$: If $n = p_1^{r_1} \dots p_\ell^{r_\ell}$, then

$$\phi(n) = p_1^{r_1-1}(p_1 - 1) \dots p_\ell^{r_\ell-1}(p_\ell - 1).$$

2. Find pos. integers u and v satisfying $ku - \phi(n)v = 1$, so that

$$ku \equiv 1 \pmod{\phi(n)}, \quad \text{i.e. } u \equiv k^{-1} \pmod{\phi(n)}.$$

3. Compute $b^u \pmod{n}$ by the method of successive squaring.

Then setting $x = \overline{b^u}$, we have

$$x^k = (\overline{b^u})^k \equiv_n b^{uk} \equiv_n b^{1+v\phi(n)} = b \cdot (b^{\phi(n)})^v \equiv_n b,$$

as desired.

Example: Find a solution to $x^{131} \equiv 758 \pmod{1073}$.

We have $\gcd(x, 1073) | \gcd(758, 1073) = 1 \checkmark$

1. Compute $\phi(n)$: Factor n to get $1073 = 29 \cdot 37$. So

$$\phi(1073) = (29 - 1)(37 - 1) = 1008.$$

2. Find pos. integers u and v satisfying $ku - \phi(n)v = 1$, so that

$$ku \equiv 1 \pmod{\phi(n)}, \quad \text{i.e. } u = k^{-1} \pmod{\phi(n)}:$$

Using the Eulidean algorithm, we get

$$\begin{array}{ll} 1008 = 131 * 7 + 91 & \text{so that} \\ 131 = 91 * 1 + 40 & \\ 91 = 40 * 2 + 11 & 1 = 4 + (-1)3 = 4 + (-1)(7 + (-1)4) \\ 40 = 11 * 3 + 7 & = 2 * 4 + (-1)7 \\ 11 = 7 * 1 + 4 & = 2(11 + (-1)7) + (-1)7 \\ 7 = 4 * 1 + 3 & = \dots = (-277) * 131 + 36 * 1008. \\ 4 = 3 * 1 + 1 & \end{array}$$

Another solution:

$$1 = (-277 + 1008) * 131 + (36 - 131) * 1008 = 731 * 131 - 95 * 1008.$$

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1. Compute $\phi(n)$: Factor n to get $1073 = 29 \cdot 37$. So

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$$ku \equiv 1 \pmod{\phi(n)}, \quad \text{i.e. } u = k^{-1} \pmod{\phi(n)}:$$

$$1 = 731 * 131 - 95 * 1008$$

3. Compute $b^u \pmod{n}$ by the method of successive squaring:

We have $731 = 2^9 + 2^7 + 2^6 + 2^4 + 2^3 + 2^1 + 1$. So using

| a | $\overline{758^{2^{a-1}}}$ | $\overline{758^{2^a}}$ | we have |
|-----|----------------------------|------------------------|---|
| 1 | 574564 | 509 | $\begin{aligned} 758^{731} &\equiv_{1073} 758^{2^9} * 758^{2^7} * 758^{2^6} \\ &\quad * 758^{2^4} * 758^{2^3} * 758^2 * 758 \\ &\equiv_{1073} (1011 * 712 * 663) \\ &\quad * (625 * 1011) * (509 * 758) \\ &\equiv_{1073} 749 * 951 * 615 \\ &\equiv_{1073} 905. \end{aligned}$ |
| 2 | 259081 | 488 | |
| 3 | 238144 | 1011 | |
| 4 | 1022121 | 625 | |
| 5 | 390625 | 53 | |
| 6 | 2809 | 663 | |
| 7 | 439569 | 712 | |
| 8 | 506944 | 488 | |
| 9 | 238144 | 1011 | |

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1. Compute $\phi(n)$: Factor n to get $1073 = 29 \cdot 37$. So

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$$ku \equiv 1 \pmod{\phi(n)}, \quad \text{i.e. } u = k^{-1} \pmod{\phi(n)}:$$

$$1 = 731 * 131 - 95 * 1008$$

3. Compute $b^u \pmod{n}$ by the method of successive squaring:

$$758^{731} \equiv_{1073} 905.$$

Then setting $x = 905$, we have

$$\begin{aligned} 905^{131} &\equiv_{1073} (758^{731})^{131} = 758^{731 \cdot 131} = 758^{1+95 \cdot 1008} \\ &= 758 \cdot (758^{1008})^{95} \equiv_{1073} 758, \end{aligned}$$

as desired. So $x = 905$ is a solution to $x^{131} \equiv 758 \pmod{1073}$.

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1. Compute $\phi(n)$: If $n = p_1^{r_1} \cdots p_\ell^{r_\ell}$, then

$$\phi(n) = p_1^{r_1-1}(p_1 - 1) \cdots p_\ell^{r_\ell-1}(p_\ell - 1).$$

(By prime factorization, which is “slow”!!)

2. Find pos. integers u and v satisfying $ku - \phi(n)v = 1$, so that

$$ku \equiv 1 \pmod{\phi(n)}, \quad \text{i.e. } u = k^{-1} \pmod{\phi(n)}.$$

(By Euclidean algorithm, which is “fast”.)

3. Compute $b^u \pmod{n}$ by the method of successive squaring.

(By method of successive squaring, which is “fast”.)

Then setting $x = \overline{b^u}$, we have

$$x^k = (\overline{b^u})^k \equiv_n b^{uk} \equiv_n b^{1+v\phi(n)} = b \cdot (b^{\phi(n)})^v \equiv_n b,$$

as desired.

How computationally difficult for large n ?

Punchline: If you know the prime factorization of n , this computation is fast (“polynomial time”); if you don’t, this computation is slow (for now—see “P versus NP”).