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2. For A a number system, let $a \sim b$ if $a < b$. **not R**
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$$(a \sim b \wedge b \sim c) \Rightarrow a \sim c$$

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An **equivalence relation** on a set A is a binary relation that is reflexive, symmetric, and transitive. **(Only #1)**

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Is \sim is an equivalence relation?

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Yes! This is an equivalence relation!

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So **no**, it is not an equivalence relation.

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Read: Why reflexivity doesn't follow from symmetry and transitivity.

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Example: Consider the equivalence relation on $A = \{a, b, c\}$ given by

$$a \sim a, \quad b \sim b, \quad c \sim c, \quad a \sim c, \quad \text{and} \quad c \sim a.$$

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Then

$$[a] = \{a, c\} = [c], \quad \text{and}$$

$$[b] = \{b\}$$

are the **two** equivalence classes in A (with respect to this relation).

(We say there are two, *not three*, since “the equivalence classes” refers to the sets themselves, not to the elements that generate them.)

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is an equivalence relation on \mathbb{Z} .

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We call any element a of a class C **representative** of C (since we can write $C = [a]$ for any $a \in C$).

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So $\mathbb{Z} = [0] \sqcup [1] \sqcup [2] \sqcup [3] \sqcup [4]$.

Ok, so what **are** numbers, anyway?

Recall from the homework, the Zermelo-Fraenkel axioms of set theory, which tells us how to compare sets, put sets in other sets, and to take unions, intersection, and power sets of sets. ✓

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Let $0 = \emptyset$.

Given n , define the **successor** to n as $S(n) = n \cup \{n\}$.

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Let \mathbb{N} be the set of all sets generated by 0 and S .

For example,

$$1 = 0 \cup \{0\} = \emptyset \cup \{\emptyset\} = \{\emptyset\},$$

$$2 =$$

$$3 =$$

Ok, so what **are** numbers, anyway?

Recall from the homework, the Zermelo-Fraenkel axioms of set theory, which tells us how to compare sets, put sets in other sets, and to take unions, intersection, and power sets of sets. ✓

Set theoretic definition of the natural numbers. ($\mathbb{Z}_{\geq 0}$)

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Think: Given $n, m \in \mathbb{N}$, are $n \cup m$ and/or $n \cap m$ elements of \mathbb{N} ? If so, what elements are they?

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Addition: Define $+$: $\mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ by, for all $a, b \in \mathbb{N}$,

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$$a + 1 = a + S(0) = S(a + 0) = S(a);$$

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$$a + 1 = a + S(0) = S(a + 0) = S(a);$$

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Check that $a + 3 = S(S(S(a))) = S^3(a)$.

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Check that $a + 3 = S(S(S(a))) = S^3(a)$. **Think:** $a + b = S^b(a)$.

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Check that $a + 3 = a + a + a$.

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Check that $a + 3 = a + a + a$.

Think: $a \cdot b = S^{ab}(0)$.

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Think: $a \cdot b = S^{ab}(0)$.

Properties:

1. Addition is commutative, i.e. $a + b = b + a$ for all $a, b \in \mathbb{N}$.
2. Addition is associative, i.e. $a + (b + c) = (a + b) + c$ for all $a, b, c \in \mathbb{N}$.
3. Multiplication is commutative, i.e. $a \cdot b = b \cdot a$ for all $a, b \in \mathbb{N}$.
4. Multiplication is associative, i.e. $a \cdot (b \cdot c) = (a \cdot b) \cdot c$ for all $a, b, c \in \mathbb{N}$.
5. Multiplication is distributive across addition, i.e.
 $a \cdot (b + c) = (a \cdot b) + (a \cdot c)$ for all $a, b, c \in \mathbb{N}$.

(These all follow from the definitions, but we'll skip proofs for the sake of time.)

Peano axioms

The natural numbers \mathbb{N} are defined by the following axioms.

1. We have $0 \in \mathbb{N}$. (technically, $0 = \emptyset$)
2. There exists an a **successor** function $S : \mathbb{N} \rightarrow \mathbb{N}$ (namely, if $n \in \mathbb{Z}$, then $S(n) \in \mathbb{N}$) satisfying
 - (i) $0 \notin S(\mathbb{N})$;
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- (v) If $a \leq b$ then $ac \leq bc$.

(These all follow from the axioms, but we'll skip proofs for the sake of time.)

Integers

Now that we have \mathbb{N} , we can define \mathbb{Z} by formally letting

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We can also extend our definitions of $+$ and \cdot to \mathbb{Z} .

Properties:

1. Addition and multiplication still satisfy commutativity, associativity, and distributivity.
2. We still have $a + 0 = a$ (**additive identity**) and $a \cdot 1 = a$ (**multiplicative identity**) for all $a \in \mathbb{Z}$.
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We call any number system that has an addition and multiplication that satisfy all these properties a **ring** (modern algebra).

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We can also extend our definition of order to \mathbb{Z} . The only modification is:

- (i) For all $a, b \in \mathbb{N}$, we have $a \leq b$ or $b \leq a$.
- (ii) If $a \leq b$ and $b \leq a$, then $a = b$.
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- (v) If $a \leq b$ and $c \in \mathbb{N}$, then $ac \leq bc$.

These properties make \mathbb{Z} an **ordered ring**.

Rational numbers

Let

$$Q = \mathbb{Z} \times (\mathbb{Z} - \{0\}),$$

and define an equivalence relation on Q by

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Under this equivalence relation, write

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This makes \mathbb{Q} a **field** (again, modern algebra).

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