Functions

Let $f: X \to Y$ be a function. Recall, the image of f is $f(X) = \{y \in Y \mid f(x) = y \text{ for some } x \in X\}.$ Further, f is...

• injective if at most one $x \in X$ maps to each $y \in Y$, i.e.

if
$$f(x_1) = f(x_2)$$
 then $x_1 = x_2$.
Ex. $f : \mathbb{R}_{>0} \to \mathbb{R}$ defined by $x \mapsto x^2$.
Non-ex. $f : \mathbb{R} \to \mathbb{R}$ defined by $x \mapsto x^2$.

• surjective if every $y \in Y$ gets mapped to, i.e.

for all $y \in Y$, there exists $x \in X$ such that f(x) = y. Ex. $f : \mathbb{R} \to \mathbb{R}_{\geq 0}$ defined by $x \mapsto x^2$. Non-ex. $f : \mathbb{R} \to \mathbb{R}$ defined by $x \mapsto x^2$.

• bijective if it's both injective and surjective.

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 defined by $x \mapsto x^2$.
Non-ex. $f : \mathbb{R} \to \mathbb{R}$ defined by $x \mapsto x^2$.

Injective, not surjective:



Surjective, not injective:



Injective, not surjective:

Surjective, not injective:



To show $f: X \to Y$ is injective: Start: Suppose $f(x_1) = f(x_2)...$ Goal: Show $x_1 = x_2$. To show $f: X \to Y$ is surjective: Start: Let $y \in Y...$ Goal: Find $x \in X$ such that f(x) = y.

$$f^{-1}(y) = \{x \in X \mid f(x) = y\} \subset X,$$

called the inverse image or preimage of y. Note that this is abusing notation: we write f^{-1} whether or not f is invertible; and $f^{-1}(y)$ is a *set*, not an *element*.

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Thm. For nonempty sets X and Y, a function $f: X \to Y$ is invertible if and only if it is bijective.

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Thm. For nonempty sets X and Y, a function $f: X \to Y$ is invertible if and only if it is bijective.

Proof.

Suppose f is bijective. Since f is surjective, for all $y \in Y$, we have $|f^{-1}(y)| \ge 1$. And since f is injective, for any $x_1, x_2 \in f^{-1}(y)$, we have $x_1 = x_2$. So $|f^{-1}(y)| \le 1$. Therefore, for all $y \in Y$, we have $|f^{-1}(y)| = 1$, so that f is invertible.

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Now suppose f is invertible. Thus for all $y \in Y$, we have $|f^{-1}(y)| = 1$. Therefore, for all $y \in Y$, $f(y) \neq \emptyset$, so that f is surjective. And for all $y \in Y$, since $f^{-1}(y)$ has exactly one element, it has at most one element. So f is injective. Therefore, f is bijective.

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Example: The sets $\mathbb{Z}_{\geqslant 0}$ and $\mathbb{Z}_{>0}$ have the same cardinality since

$$\begin{array}{rccc} f: \mathbb{Z}_{>0} & \to & \mathbb{Z}_{\geq 0} \\ x & \mapsto & x-1 \end{array}$$

is a bijective map.

A set is countable if it is either finite or the same cardinality as the natural numbers $(\mathbb{N} = \mathbb{Z}_{>0})$. If a set A is not finite but is countable, we say A is "countably infinite" and write $|A| = \aleph_0$ (pronounced "aleph naught" or "aleph null").

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 Z is countably infinite: Not finite, and f: Z_{>0} → Z by x ↦ (-1)^x[x/2] is a bijection.
 Z_{>0}: ... 9 7 5 3 1 2 4 6 8 ... ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ Z: ... -4 -3 -2 -1 0 1 2 3 4 ... More on this last example, $|\mathbb{Z}| = \aleph_0$:

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This at least gives us a "list" of integers,

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If I know that every integer appears on this list somewhere, then I know that the integers are countable. (Ok answer)
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$$f: \mathbb{Z}_{>0} \to \mathbb{Z} \qquad x \mapsto \begin{cases} x/2 & \text{if } x \text{ is even,} \\ -(x-1)/2 & \text{if } x \text{ is odd.} \end{cases}$$
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To be even more sophisticated, we used the floor function to get a closed form answer:

$$\begin{cases} x/2 & \text{if } x \text{ is even,} \\ -(x-1)/2 & \text{if } x \text{ is odd} \end{cases} = (-1)^x \lfloor x/2 \rfloor,$$
 for $x \in \mathbb{Z}_{>0}$. (Best answer)

Recall that |A| = |B| if and only if there is a bijection $f : A \rightarrow B$.

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Examples: For each of the following, show that the set is countably infinite. (Define a bijective function to something that we know to be countably infinite if it's not too hard; otherwise, explain how to make the list.)

- 1. The set of negative integers $(\mathbb{Z}_{<0})$.
- 2. The set of integers less than 100 ($\mathbb{Z}_{<100}$).
- 3. The set of integers that are integer multiples of 3 ($3\mathbb{Z}$).
- 4. The set of integers that are not integer multiples of 3 $(\mathbb{Z} 3\mathbb{Z})$.

Claim: $\mathbb{Q}_{>0}$ is countably infinite.

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		1	2	3	4	5	
	1	1/1	2/1	3/1	4/1	5/1	
	2	1/2	2/2	3/2	4/2	5/2	
	3	1/3	2/3	3/3	4/3	5/3	
	4	1/4	2/4	3/4	4/4	5/4	
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Then use the same alternating map that we did for $\mathbb{Z}_{>0} \to \mathbb{Z}$ to build a bijection $\mathbb{Q}_{>0} \to \mathbb{Q}$, yielding a bijection $\mathbb{Z}_{>0} \to \mathbb{Q}_{>0} \to \mathbb{Q}$.

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Proof outline: We will prove this by contradiction. Suppose that the set [0,1) is countable, so that the real numbers in [0,1) can be listed. Take one such list.

Goal: Show that the list isn't complete. Namely, algorithmically produce an element of $\left[0,1\right)$ that isn't on any fixed list.

For example:

- 1. 0.001240191057...
- 2. 0.123451234512...
- **3**. 0.333333333333...
- 4. 0.50000000000...
- 5. 0.121212121212...
- **6**. 0.555555555555...
- 7. 0.141592653589...
- 8. 0.00185000000...
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- **1**. 0.**0**01240191057...
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Build a new number as follows:

- i. If the highlighted digit of the *i*th number is a 0, then make the corresponding digit of the new number a 1.
- ii. If the highlighted digit of the *i*th number is not a 0, then make the corresponding digit of the new number a 0.

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Example:

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In this way, this new number differs from every item in the list in at least one digit!

Proof outline: We will prove this by contradiction. Suppose that the set [0,1) is countable, so that the real numbers in [0,1) can be listed. Show that any list isn't complete. Namely, algorithmically produce an element of [0,1) that isn't on any fixed list.

Proof. For $x \in [0, 1)$, denote the *i*th digit of x by x[i]. Note that for $x, y \in [0, 1)$, we have x = y if and only if x[i] = y[i] for all $i \in \mathbb{Z}_{>0}$.

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Recall, the power set of as set A is the set of subsets of A, given by

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Theorem. $|A| < |\mathcal{P}(A)|$.

Outline:

- 1. Show $|A| \leq |\mathcal{P}(A)|$ by showing an injective map *exists* (give one example).
- 2. Show |A| ≠ |P(A)| by showing that any map (not just the example from before) cannot be surjective.
 Hint. For any f : A → P(A), show f is not surjective; i.e. construct a set S ⊆ A such that S ≠ f(a) for all a ∈ A.

(Need for care: 1 is a "there exists" statement; 2 is a "for all" statement.)

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Let \sim be the relation on \mathbb{Z} given by $a \sim b$ if a < b. (Note that we use language like in definitions, where "if" actually means "if and only if".)

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A binary relation on a set A is... (R) reflexive if $a \sim a$ for all $a \in A$;

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- 2. For A a number system, let $a \sim b$ if a < b. not R
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- 1. For A a number system, let $a \sim b$ if a = b. R, S, T
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An equivalence relation on a set A is a binary relation that is reflexive, symmetric, and transitive.
More examples of (binary) relations:

- 1. For A a number system, let $a \sim b$ if a = b. R, S, T
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An equivalence relation on a set A is a binary relation that is reflexive, symmetric, and transitive. (Only #1)

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Is \sim is an equivalence relation?

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Check: we have $a \equiv b \pmod{n}$ if and only if a - b = kn for some $k \in \mathbb{Z}$.

"
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Is \sim is an equivalence relation?

Check: we have $a \equiv b \pmod{n}$ if and only if a - b = kn for some $k \in \mathbb{Z}$.

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$$a-c$$

"
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Check: we have $a \equiv b \pmod{n}$ if and only if a - b = kn for some $k \in \mathbb{Z}$.

$$a - c = (a - b) + (b - c)$$

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$$a \sim b$$
 if $a \equiv b \pmod{n}$."

Is \sim is an equivalence relation?

Check: we have $a \equiv b \pmod{n}$ if and only if a - b = kn for some $k \in \mathbb{Z}$.

$$a - c = (a - b) + (b - c) = kn + \ell n$$

$$a \sim b \quad \text{if } a \equiv b \pmod{n}.$$

Is \sim is an equivalence relation?

Check: we have $a \equiv b \pmod{n}$ if and only if a - b = kn for some $k \in \mathbb{Z}$.

$$a - c = (a - b) + (b - c) = kn + \ell n = (k + \ell)n.\checkmark$$

$$a \sim b \quad \text{if } a \equiv b \pmod{n}.$$

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Check: we have $a \equiv b \pmod{n}$ if and only if a - b = kn for some $k \in \mathbb{Z}$.

reflexivity: $a - a = 0 = 0 \cdot n \checkmark$ symmetry: If a - b = kn, then b - a = -kn = (-k)n. \checkmark transitivity: If a - b = kn and $b - c = \ell n$, then

$$a - c = (a - b) + (b - c) = kn + \ell n = (k + \ell)n.\checkmark$$

Yes! This is an equivalence relation!

Let A be a set. Consider the relation on $\mathcal{P}(A)$ by $S\sim T \qquad \text{if} \qquad S\subseteq T$

Is \sim is an equivalence relation?

 $S \sim T$ if $S \subseteq T$

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Check: This is reflexive and transitive, but not symmetric. So no, it is not an equivalence relation.

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ls

$$S \sim T$$
 if $S \subseteq T$ or $S \subseteq T$

an equivalence relation on $\mathcal{P}(A)$?

 $S \sim T$ if $S \subseteq T$

Is \sim is an equivalence relation?

Check: This is reflexive and transitive, but not symmetric. So no, it is not an equivalence relation.

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$$S \sim T \qquad \text{ if } \qquad S \subseteq T \text{ or } S \subseteq T$$

an equivalence relation on $\mathcal{P}(A)$?

Check: This is reflexive and symmetric, but not transitive. So still no, it is not an equivalence relation.

 $S \sim T$ if $S \subseteq T$

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an equivalence relation on $\mathcal{P}(A)$?

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ls

$$S \sim T$$
 if $|S| = |T|$ an equivalence relation on $\mathcal{P}(A)$?

 $S \sim T$ if $S \subseteq T$

Is \sim is an equivalence relation?

Check: This is reflexive and transitive, but not symmetric. So no, it is not an equivalence relation.

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$$S \sim T \qquad \text{ if } \qquad S \subseteq T \text{ or } S \subseteq T$$

an equivalence relation on $\mathcal{P}(A)$?

Check: This is reflexive and symmetric, but not transitive. So still no, it is not an equivalence relation.

ls

$$S \sim T$$
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an equivalence relation on $\mathcal{P}(A)$?

Read: Why reflexivity doesn't follow from symmetry and transitivity.

Example: Consider the equivalence relation on $A = \{a, b, c\}$ given by

 $a \sim a, \quad b \sim b, \quad c \sim c, \quad a \sim c, \quad \text{and} \quad c \sim a.$

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[a]

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$$[a] = \{a, c\}$$

Example: Consider the equivalence relation on $A = \{a, b, c\}$ given by

 $a \sim a$, $b \sim b$, $c \sim c$, $a \sim c$, and $c \sim a$.

$$[a] = \{a, c\} = [c]$$

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Then

$$[a] = \{a, c\} = [c],$$
 and
 $[b] = \{b\}$

are the two equivalence classes in A (with respect to this relation).

(We say there are two, *not three*, since "the equivalence classes" refers to the sets themselves, not to the elements that generate them.)

Let \sim be an equivalence relation on a set A, and let $a \in A$. The set of all elements $b \in A$ such that $a \sim b$ is called the equivalence class of a, denoted by [a]. Example: We showed that " $a \sim b$ if $a \equiv b \pmod{5}$ "

is an equivalence relation on \mathbb{Z} .

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 $[0] = \{5n \mid n \in \mathbb{Z}\} = 5\mathbb{Z} \qquad [1] = \{5n+1 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 1$

Let \sim be an equivalence relation on a set A, and let $a \in A$. The set of all elements $b \in A$ such that $a \sim b$ is called the equivalence class of a, denoted by [a]. Example: We showed that " $a \sim b$ if $a \equiv b \pmod{5}$ " is an equivalence relation on \mathbb{Z} . Then $[0] = \{5n \mid n \in \mathbb{Z}\} = 5\mathbb{Z}$ $[1] = \{5n + 1 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 1$ [2]

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 $[0] = \{5n \mid n \in \mathbb{Z}\} = 5\mathbb{Z} \qquad [1] = \{5n+1 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 1$ $[2] = \{5n+2 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 2 \qquad [3] = \{5n+3 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 3$

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 $[5] = \{5n + 5 \mid n \in \mathbb{Z}\} = \{5m \mid m \in \mathbb{Z}\} = [0] = [-5] = [10] = \cdots$

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In general, if $x \in [y]$, that means $y \sim x$.

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In general, if $x \in [y]$, that means $y \sim x$. So $x \sim y$.

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In general, if $x \in [y]$, that means $y \sim x$. So $x \sim y$. So $y \in [x]$.

Let \sim be an equivalence relation on a set A, and let $a \in A$. The set of all elements $b \in A$ such that $a \sim b$ is called the equivalence class of a, denoted by [a]. Example: We showed that " $a \sim b$ if $a \equiv b \pmod{5}$ " is an equivalence relation on \mathbb{Z} . Then $[0] = \{5n \mid n \in \mathbb{Z}\} = 5\mathbb{Z} \qquad [1] = \{5n+1 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 1$ $[2] = \{5n+2 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 2 \qquad [3] = \{5n+3 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 3$ $[4] = \{5n + 4 \mid n \in \mathbb{Z}\} = 5\mathbb{Z} + 4$ $[5] = \{5n + 5 \mid n \in \mathbb{Z}\} = \{5m \mid m \in \mathbb{Z}\} = [0] = [-5] = [10] = \cdots$ $[6] = \{5n + 6 \mid n \in \mathbb{Z}\} = \{5m + 1 \mid m \in \mathbb{Z}\} = [1] = [-4] = [11] = \cdots$ In general, if $x \in [y]$, that means $y \sim x$. So $x \sim y$. So $y \in [x]$. Claim: $x \in [y]$ if and only if [x] = [y].

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We call any element a of a class C representative of C (since we can write C = [a] for any $a \in C$).

Theorem. The equivalence classes of \boldsymbol{A} partition \boldsymbol{A} into subsets, meaning

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 $[a] \subseteq A$ for all $a \in A$;

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So
$$\mathbb{Z} = [0] \sqcup [1] \sqcup [2] \sqcup [3] \sqcup [4]$$
.