PROOF LAB II: INDUCTION – HINTS AND TIPS.

Outlining your proof:

- 1. Define P(n).
- Rewrite the claim in the best way for you to use induction.
- 2. Compute base case(s). Usually P(0) or P(1). You might have to do more than one base case to get the induction step to work. 3. Explicitly state your goal.
- "Assume P(n) and prove P(n+1), which is..."
- 4. Do inductive step. This is usually where the real work happens.
- 5. State your conclusion.

Rewrite your proof: Your final draft should read something like

"We will prove this claim by induction on n. First, for n = a," we have [COMPUTATION]. Next, fix $n \ge a$ and assume [STATE THE CLAIM] for that value of n. Then [USE THAT ASSUMPTION TO SHOW CLAIM FOR n+1^{**}. Thus, by induction, the claim holds for all $n \ge a$."

Note: In my rewrite, I didn't ever say anything about P(n). For example, compare the following two writeups of the same inductive proof of $n < 2^n$ for all $n \in \mathbb{Z}_{>0}$.

Proof by induction (first draft – don't turn this in).

Define P(n): P(n) is " $n < 2^n$ ". **Base case:** The least value of n is 0, so the base case is P(0): $0 < 1 = 2^0$. \checkmark **Goal:** Assume P(n) and show P(n+1), which is $n+1 < 2^{n+1}.$ P(n+1):**Inductive step:** (Assume P(n) and show P(n+1)) Fix $n \ge 0$ and assume $n < 2^n$ (this is the IH). Then since $n \ge 0$, $n+1 \stackrel{\text{IH}}{<} 2^n + 1 \le 2^n + 2^n = 2(2^n) = 2^{n+1}.$ \checkmark **Conclusion:** So since P(0) is true, and P(n) implies P(n+1), we have P(k) is true for all $k \in \mathbb{Z}_{>0}.$

Proof (final draft). For n = 0, we have

 $0 < 1 = 2^0$,

as desired. Now, fix $n \ge 0$ and assume $n < 2^n$ (for that n). Then since $n \ge 0$, we have $n+1 < 2^n + 1 \le 2^n + 2^n = 2(2^n) = 2^{n+1}.$ Thus, the claim holds for all $n \ge 0$ by induction.

^{*}where a is something like 0 or 1 or 5—whatever the lower end of the domain of the problem is.

^{**}Don't forget to point out where you use the "Induction Hypothesis".

Problems. Prove each of the following using proof by induction. Assume throughout that n is an integer.

I. (\star) If $x_1, x_2, ..., x_n$ are odd integers, then their product,

$$\prod_{i=1}^{n} x_i \qquad \text{is also odd.}$$

[Note that $x_1, x_2, ..., x_n$ is an arbitrary list of odd numbers (you don't get to choose what these numbers are!).]

Hint: Since x_i is odd, for each i, we can write $x_i = 2k_i + 1$ for some $k_i \in \mathbb{Z}$. Start with $\prod_{i=1}^{1} x_i$. Then for $n \ge 1$, you have $\prod_{i=1}^{n+1} x_i = (\prod_{i=1}^{n} x_i) x_{i+1}$.

II. (*) For all odd natural numbers n, $n^2 - 1$ is divisible by 8. [Warning: Do not induct on n. Instead, start by writing what it means for n to be odd.]

Hint: Since n is odd, we can write n = 2k + 1 for some $k \in \mathbb{Z}$. As implied above, you don't want to induct on n; you want to induct on k, starting with k = 0 (careful!! if n = 1, then k = 0). Then $n^2 - 1$ is divisible by 8 is equivalent to $n^2 - 1 = 8\ell$ for some $\ell \in \mathbb{Z}$. Plug in n = 2k + 1 into that equation and simplify. Then move on to n = 2(k + 1) + 1.

III. $(\star\star)$ For all $n \ge 1$ and $0 \le x \le \pi$, we have $\sin(nx) \le n \sin(x)$. [You may use the angle addition formula for $\sin(x)$.]

Hint: Expand sin(nx+x) using the angle addition formula. Then you can use the fact that $cos(x) \leq 1$ for all x.

IV. $(\star\star)$ Let $a_n = 2^{2^n} + 1$. Then for all $n \ge 2$, the last digit of a_n is 7. [†] [Hint: Rephrase "last digit is 7" in terms of remainders.]

Hint: "The last digit of a_n is 7" is equivalent to $a_n = 10k + 7$ for some $k \in \mathbb{Z}_{\geq 0}$.

V. $(\star \star \star)$ The Fibonacci Numbers F_1, F_2, \ldots , are defined by the recursive rule

 $F_1 = 1$ $F_2 = 1$ and $F_{n+1} = F_n + F_{n-1}$ for all $n \ge 3$.

For all $n \ge 1$, F_{5n} is divisible by 5.

[Hint: For the induction step, let r be the remainder of F_{5n-1} when divided by 5.]

Hint: Working backwards, you have $F_{5n+5} = F_{5n+4} + F_{5n+3}$, $F_{5n+4} = F_{5n+3} + F_{5n+2}$, and so on. You will assume something about F_{5n} , so that's how far back you'll need to take that computation.

[†]Fermat conjectured that for $n \ge 1$, a_n is always prime. Indeed, $a_1 = 5, a_2 = 17, a_3 = 257, a_4 = 65, 537$ are all prime. However, the conjecture fails for many large values of n (Fermat had a hard time checking because a_i gets really large quickly).

VI. $(\star \star \star)$ If u and v are differentiable functions of x, then

$$\frac{d^n}{dx^n}uv = \sum_{k=0}^n \binom{n}{k} u_k v_{n-k},$$

where $u_i = \frac{d^i}{dx^i}u$ and $v_i = \frac{d^i}{dx^i}v$. [You will need to prove a lemma that shows $\binom{r}{s} + \binom{r}{s+1} = \binom{r+1}{s+1}$ for all $r, s \in \mathbb{Z}_{\geq 0}$, to be proven directly.]

Hint: Just try out the first few steps and see what happens. Namely,

$$\begin{aligned} \frac{d}{dx}uv &= \left(\frac{d}{dx}u\right)v + u\left(\frac{d}{dx}v\right) = u_1v_0 + u_0v_1;\\ \frac{d^2}{dx^2}uv &= \frac{d}{dx}\left(u_1v_0 + u_0v_1\right)\right)\\ &= \left(\frac{d}{dx}u_1\right)v_0 + u_1\left(\frac{d}{dx}v_0\right) + \left(\frac{d}{dx}u_0\right)v_1 + u_0\left(\frac{d}{dx}v_1\right)\right)\\ &= u_2v_0 + u_1v_1 + u_1v_1 + u_0v_2\\ &= u_2v_0 + 2u_1v_1 + u_0v_2 = \sum_{i=0}^2 \binom{2}{i}u_iv_{2-i}; \quad \text{and}\\ \frac{d^3}{dx^3}uv &= \frac{d}{dx}\left(\frac{d^2}{dx^2}uv\right) = \frac{d}{dx}\left(\sum_{i=0}^2 \binom{2}{i}u_iv_{2-i}\right)\\ &= \sum_{i=0}^2 \binom{2}{i}\frac{d}{dx}(u_iv_{2-i}) = \sum_{i=0}^2 \binom{2}{i}(u_{i+1}v_{2-i} + u_iv_{2-i+1})\\ &= \sum_{i=0}^2 \binom{2}{i}u_{i+1}v_{2-i} + \sum_{i=0}^2 \binom{2}{i}u_iv_{2-i+1}\\ &= (u_1v_2 + 2u_2v_1 + u_3v_0) + (u_0v_3 + 2u_1v_2 + u_2v_1)\\ &= u_0v_3 + 3u_1v_2 + 3u_2v_1 + u_3v_0.\end{aligned}$$

Basically, this works exactly the same as the binomial theorem! To show the lemma $\binom{r}{s} + \binom{r}{s+1}$, just plug in $\binom{a}{b} = \frac{a!}{b!(a-b)!}$, and simplify.