An Algebraic Approach to Voting Theory

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So who should be president?

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Method 1: Vote for your first choice.

$$\begin{array}{c|cc} A & B & C \\ \hline 13 & 19 & 22 \end{array}$$

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Method 2: Tell us your full ranking, and we'll pair them off.

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A > B	B > C	A > C
30	32	32
A < B	B < C	A < C

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A > B	B > C	A > C
30	32	32
A < B	B < C	A < C
24	22	22
A > B,	B > C,	A > C.

So who should be president?

Method 1: Vote for your first choice.

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A > B	B > C	A > C
30	32	32
A < B	B < C	A < C
24	22	22
A > B,	B > C,	A > C.

Alice wins!

So who should be president?

Method 3: Tell us your full ranking, and we'll run an instant run off.

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Round 1:ABC131922

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Round 1:ABC131922

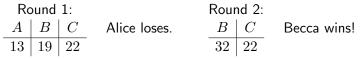
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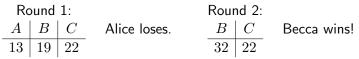
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$$\frac{A \mid B \mid C}{13 + 36t \mid 19 + 18t \mid 22 + 0t} \quad (0 \le t \le 1)$$

So who should be president?

		Α	В	C	
	$(0 \leqslant t \leqslant 1)$	13 + 36t	19 + 18t	22 + 0t	
	t = 0:	13	19	22	
	t = 0.25:	22	23.5	22	
	t = 0.5 :	31	28	22	
	t = 0.75:	40	32.5	22	
	t = 1:	49	37	22	
C > B >	$A \qquad B > C$	$> A \mid B > A$	A > C	$\underline{A} > \underline{B} > \underline{C}$	
	1/6	1/4	1/3		1

$1^{st}$ :	B	G	G	N
$2^{nd}$ :	G	N	B	G
$3^{rd}$ :	N	B	N	В
	2.902mil	1.421mil	1.481mil	0.118mil

So who should have been president?

$1^{st}$ :	B	G	G	N
$2^{nd}$ :	G	N	B	G
$3^{rd}$ :	N	B	N	В
	2.902mil	1.421mil	1.481mil	0.118mil

So who should have been president?

Method 1: Vote for your first choice.

B	G	N	Bush won
2.90mil	2.90mil	0.11mil	Dusit won

$1^{st}$ :	B	G	G	N
$2^{nd}$ :	G	N	B	G
$3^{rd}$ :	N	В	N	В
	2.902mil	1.421mil	1.481mil	0.118mil

So who should have been president?

Method 1: Vote for your first choice.

Method 2: Tell us your full ranking, and we'll pair them off.

B > G	G > N	B > N
2.90mil	5.80mil	4.32mil
B < G	G < N	B < N
3.02mil	0.12mil	1.60mil
$\overline{G > B},$	G > N,	B > N.

Gore wins.

$1^{st}$ :	B	G	G	N
$2^{nd}$ :	G	N	B	G
$3^{rd}$ :	N	B	N	В
	2.902mil	1.421mil	1.481mil	0.118mil

So who should have been president?

Method 3: Tell us your full ranking, and we'll run an instant run off.

Round 1:

B	G	N
2.90mil	2.90mil	0.12mil

$1^{st}$ :	B	G	G	N
$2^{nd}$ :	G	N	B	G
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$1^{st}$ :	B	G	G	N
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	2.902mil	1.421mil	1.481mil	0.118mil

So who should have been president?

Method 3: Tell us your full ranking, and we'll run an instant run off.

	Round 1:		Round 2:		
B	G	N	Nader loses.	B	G
2.90mil	2.90mil	0.12mil		2.90mil	3.02mil

$1^{st}$ :	B	G	G	N
$2^{nd}$ :	G	N	B	G
$3^{rd}$ :	N	B	N	В
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$2^{nd}$ :	G	N	B	G
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So who should have been president?

Method 3: Tell us your full ranking, and we'll run an instant run off.

Method 4: Tell us your full ranking, and we'll give points according to the candidates' positions—1 for 1st, t for 2nd, 0 for 3rd.

$1^{st}$ :	B	G	G	N
$2^{nd}$ :	G	N	B	G
$3^{rd}$ :	N	B	N	В
	2.902mil	1.421mil	1.481mil	0.118mil

So who should have been president?

		A	В	C
	$(0 \leqslant t \leqslant 1)$	2.902 + 1.421 t	2.902 + 3.020 t	0.118 + 1.481 t
	t = 0:	2.90mil	2.90mil	0.19mil
	t = 0.25 :	3.26mil	3.66mil	0.49mil
	t = 0.5 :	3.61mil	4.41mil	0.86mil
	t = 0.75:	3.97mil	5.17mil	1.23mil
	t = 1:	4.32mil	5.92mil	1.60mil
+	B > G > N		G > B > N	
ι (	l .000	335		



Jean-Charles, Chevalier de Borda 1733–1799 Mariner and scientist. 1770: formulated a ranked voting system, the "Borda count". Used by the French Academy of Sciences, until Napolean.

#### Nicolas de Caritat, Marquis de Condorcet 1743–1794

Philosopher and mathematician.

In 1785, wrote an essay on probability of decisions made on a majority vote, describing likelihood of good jury outcomes; and Condorcet's paradox, which shows that majority preferences can become intransitive with three or more options





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Professor of Mathematics and Economics.

1999: Used geometric methods to model voting data as vector spaces, and decompose them based on how they affect various tallying methods.



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Kernel: Doesn't affect any fair voting system.

Borda: Influences both point-based and pairwise systems.

Condorcet: Introduces Condorcet paradox.

Reversal: Influences point-based systems, but not pairwise.

	Kernel	Borda	Condorcet	Reversal
		$\mathbf{b}_{\mathbf{A}}$ $\mathbf{b}_{\mathbf{B}}$ $\mathbf{b}_{\mathbf{C}}$		$\mathbf{r}_{\mathbf{A}}$ $\mathbf{r}_{\mathbf{B}}$ $\mathbf{r}_{\mathbf{C}}$
ABC	$\langle 1 \rangle$	$\begin{pmatrix} 1 \\ -1 \end{pmatrix} \begin{pmatrix} 0 \\ -1 \end{pmatrix}$	$\begin{pmatrix} 1 \end{pmatrix}$	$\begin{pmatrix} 1 \\ -2 \\ -1 \end{pmatrix}$
ACB	1	1 -1 0	-1	$1 \ 1 \ -2$
BAC	1	0 1 -1	-1	-2 1 1
BCA	1		1	1   1   -2
CAB	1	0 -1 1	1	-2  1 1
CBA		$\left  \begin{array}{c} -1 \end{array} \right  \left  \begin{array}{c} 0 \end{array} \right  \left  \begin{array}{c} 1 \end{array} \right $	$\left(-1\right)$	1/-2/1/1/

$1^{st}$ :	A	A	B	B	C	C	Points			A > B	B > C	C > A
$2^{nd}$ :	B	C	A	C	A	B	w	t = (1, t, t)	0) (0	VS.	VS.	VS.
$3^{rd}$ :	C	B		A	B	A	A	B	C	A < B	B < C	C < A
Ker	1	1	1	1	1	1	2+2t	2 + 2t	2+2t	0	0	0
$\mathbf{b}_A$	1	1	0	-1	0	-1	2	-1	-1	4	0	-4
$\mathbf{b}_B$	0	-1	1	1	-1	0	-1	2	-1	-4	4	0
$\mathbf{b}_C$	-1	0	-1	0	1	1	-1	-1	2	0	-4	4
Cond	1	-1	-1	1	1	-1	0	0	0	2	2	2
$\mathbf{r}_A$	1	1	-2	1	-2	1	2-4t	-1+2t	-1+2t	0	0	0
$\mathbf{r}_B$	-2	1	1	1	1	-2	-1+2t	2-4t	-1+2t	0	0	0
$\mathbf{r}_C$	1	-2	1	-2	1	1	-1+2t	-1+2t	2-4t	0	0	0

Don't worry about negatives: negative votes are fixed by the kernel.

$1^{st}$ :	A	A	B	B	C	C	Points			A > B	B > C	C > A
$2^{nd}$ :	B	C	A	C	A	B	w	t = (1, t, t)	0) (0	VS.	VS.	VS.
$3^{rd}$ :	C	B		A	B	A	A	B	C	A < B	B < C	C < A
Ker	1	1	1	1	1	1	2+2t	2 + 2t	2 + 2t	0	0	0
$\mathbf{b}_A$	1	1	0	-1	0	-1	2	-1	-1	4	0	-4
$\mathbf{b}_B$	0	-1	1	1	-1	0	-1	2	-1	-4	4	0
$\mathbf{b}_C$	-1	0	-1	0	1	1	-1	-1	2	0	-4	4
Cond	1	-1	-1	1	1	-1	0	0	0	2	2	2
$\mathbf{r}_A$	1	1	-2	1	-2	1	2-4t	-1+2t	-1+2t	0	0	0
$\mathbf{r}_B$	-2	1	1	1	1	-2	-1+2t	2-4t	-1+2t	0	0	0
$\mathbf{r}_{C}$	1	-2	1	-2	1	1	-1+2t	-1+2t	2-4t	0	0	0

Don't worry about negatives: negative votes are fixed by the kernel. Note: The reversal space is trivial *precisely* when t = 1/2.

		Points	A > B	B > C	C > A	
	v	$\mathbf{v} = (1, t, 0)$	)	VS.	VS.	VS.
	A	В	C	A < B	B < C	$C\!<\!A$
Ker	2+2t	2 + 2t	2+2t	0	0	0
$\mathbf{b}_A$	2	-1	-1	4	0	-4
$\mathbf{b}_B$	-1	2	-1	-4	4	0
$b_C$	-1	-1	2	0	-4	4
Cond	0	0	0	2	2	2
$\mathbf{r}_A$	2-4t	-1 + 2t	-1 + 2t	0	0	0
$\mathbf{r}_B$	-1+2t	2-4t	-1 + 2t	0	0	0
$\mathbf{r}_{C}$	-1+2t	-1 + 2t	2-4t	0	0	0

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		Points	A > B	B > C	C > A	
	v	$\mathbf{v} = (1, t, 0)$	)	VS.	VS.	VS.
	A	В	C	$A \! < \! B$	B < C	$C\!<\!A$
Ker	2+2t	2 + 2t	2+2t	0	0	0
$\mathbf{b}_A$	2	-1	-1	4	0	-4
$\mathbf{b}_B$	-1	2	-1	-4	4	0
$b_C$	-1	-1	2	0	-4	4
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$\mathbf{r}_B$	-1+2t	2-4t	-1 + 2t	0	0	0
$\mathbf{r}_{C}$	-1+2t	-1 + 2t	2-4t	0	0	0

Don't worry about negatives: negative votes are fixed by the kernel. Note: The reversal space is trivial *precisely* when t = 1/2.

Theorem (Saari '00) Given a full ranking of n candidates, the reversal space is trivial precisely for weight

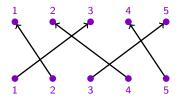
$$\mathbf{w} = \left(1, \ \frac{n-2}{n-1}, \ \dots, \ \frac{2}{n-1}, \ \frac{1}{n-1}, \ 0\right).$$

## Permutations and the symmetric group

A permutation is a bijective (one-to-one and onto) function

$$\sigma: \{1, \ldots, n\} \to \{1, \ldots, n\}.$$

Permutation diagrams:

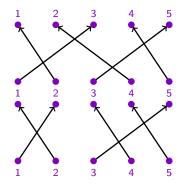


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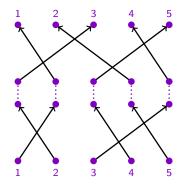
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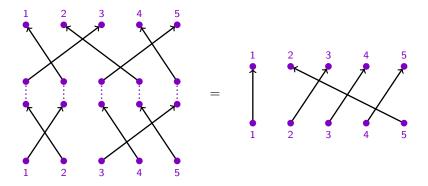
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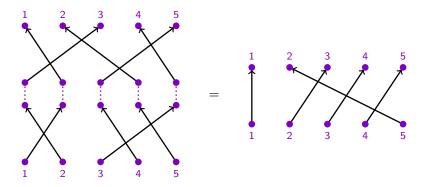
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# Permutations and the symmetric group

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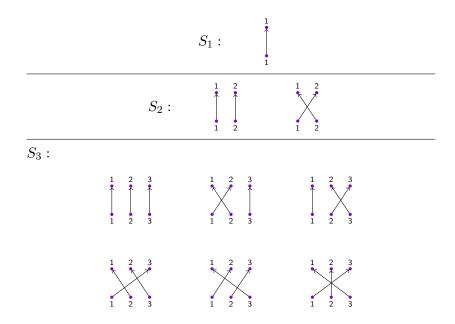
Permutation diagrams:



Permutations "multiply" by stacking and resolving.

The symmetric group  $S_n$  is the group of permutations of  $1, \ldots, n$  with multiplication given by function composition.

### Some examples:



Example: Permutation representation of the symmetric group.

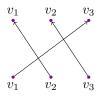


Example: Permutation representation of the symmetric group.



Pick a basis for 
$$\mathbb{Q}^3$$
:  
 $v_1 = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$   $v_2 = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}$   $v_3 = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$ 

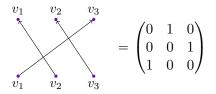
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Map each permutation to the matrix that permutes the basis vectors in the same way. (Recall: *i*th col. is image of *i*th basis vector)

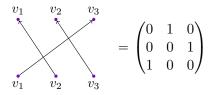
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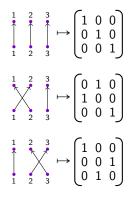
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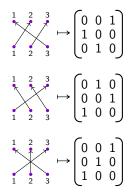


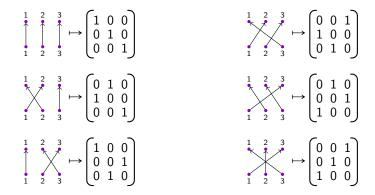
Pick a basis for  $\mathbb{Q}^3$ :  $v_1 = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$   $v_2 = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}$   $v_3 = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$ Map each permutation to the matrix that permutes the basis

vectors in the same way. (Recall: *i*th col. is image of *i*th basis vector)

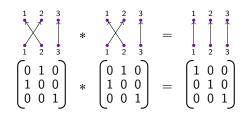
Aside: we actually have a representation of the group ring  $\mathbb{Q}S_n = \left\{\sum_{\sigma \in S_n} r_{\sigma}\sigma \mid r_{\sigma} \in \mathbb{Q}\right\}$ , with multiplication like polynomials.

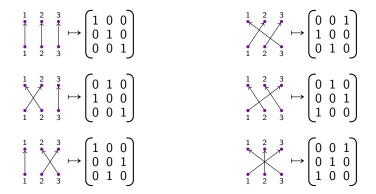




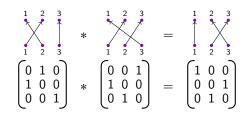


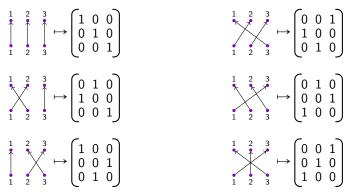
For example,





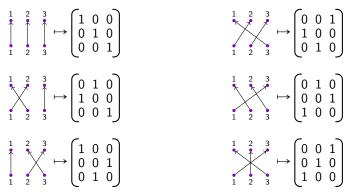
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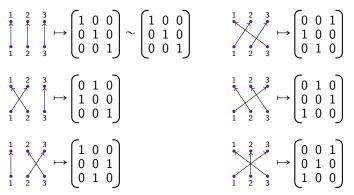
$$M(v_1 + v_2 + v_3) = v_1 + v_2 + v_3$$

for all permutation matrices M.



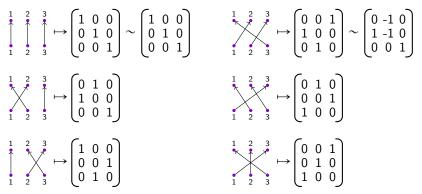
$$M(v_1 + v_2 + v_3) = v_1 + v_2 + v_3$$

$$w_1 = v_1 - v_2,$$
  $w_2 = v_2 - v_3,$   $w_3 = v_1 + v_2 + v_3$ 



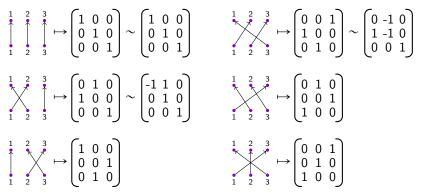
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$$M(v_1 + v_2 + v_3) = v_1 + v_2 + v_3$$

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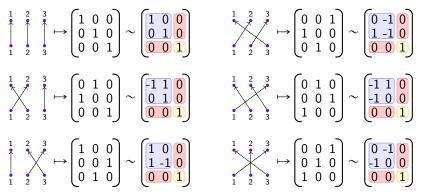
$$M(v_1 + v_2 + v_3) = v_1 + v_2 + v_3$$

$$w_1 = v_1 - v_2, \qquad w_2 = v_2 - v_3, \qquad w_3 = v_1 + v_2 + v_3$$

$$\begin{array}{c} 1 & 2 & 3 \\ 1 & 2 & 3 \\ 1 & 2 & 3 \end{array} \mapsto \left( \begin{array}{c} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{array} \right) \sim \left( \begin{array}{c} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{array} \right) \qquad \begin{array}{c} 1 & 2 & 3 \\ 1 & 2 & 3 \end{array} \mapsto \left( \begin{array}{c} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{array} \right) \sim \left( \begin{array}{c} 0 -1 & 0 \\ 1 -1 & 0 \\ 0 & 0 & 1 \end{array} \right)$$
$$\begin{array}{c} 1 & 2 & 3 \\ 1 & 2 & 3 \end{array} \mapsto \left( \begin{array}{c} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{array} \right) \sim \left( \begin{array}{c} -1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{array} \right)$$
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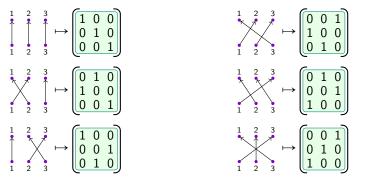
$$M(v_1 + v_2 + v_3) = v_1 + v_2 + v_3$$

$$w_1 = v_1 - v_2, \qquad w_2 = v_2 - v_3, \qquad w_3 = v_1 + v_2 + v_3$$

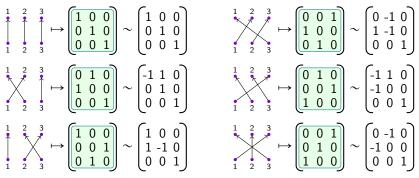


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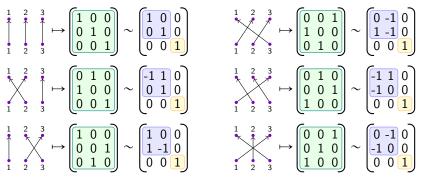
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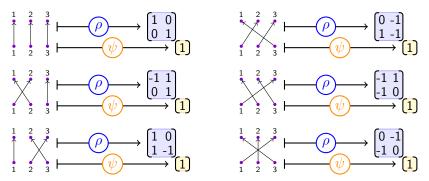
Start with the permutation representation P with basis  $\{v_1, v_2, v_3\}$ .



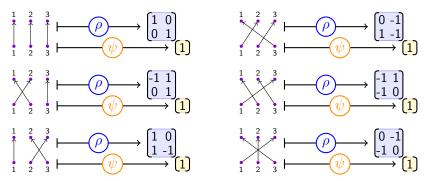
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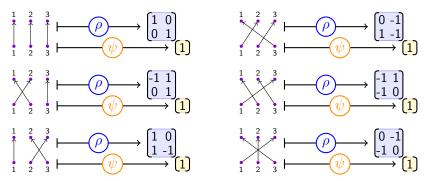


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We say P is isomorphic to the sum of two smaller representations:  $P\cong\rho\oplus\psi$ We say  $\rho$  and  $\psi$  are simple because neither has any invariant

subspaces.

### Some combinatorics.

Let n be a non-negative integer. A partition  $\lambda$  of n is a non-ordered list of positive integers which sum to n.

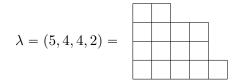
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We draw partitions as n boxes left-justified, where the parts are the number of boxes in a row (reading from the bottom):

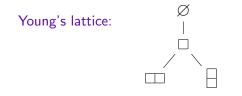


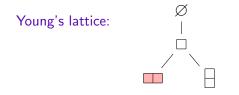
Young's lattice:

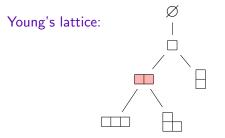
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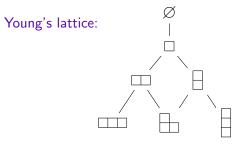


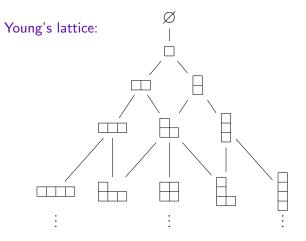


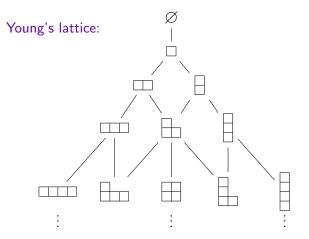




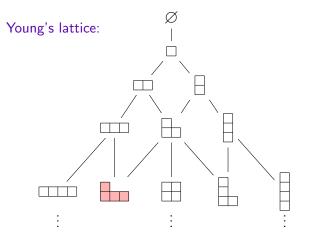




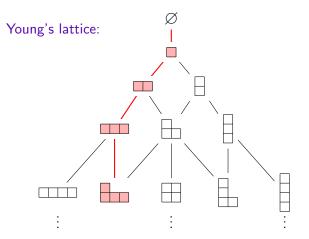




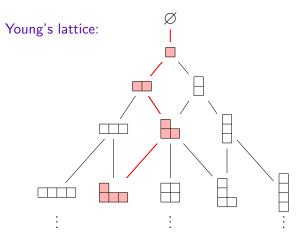
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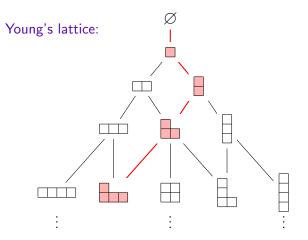
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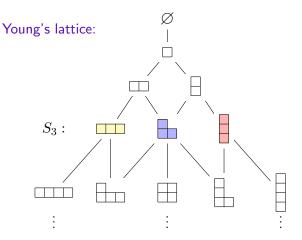
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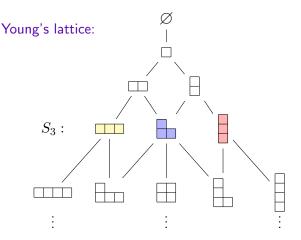
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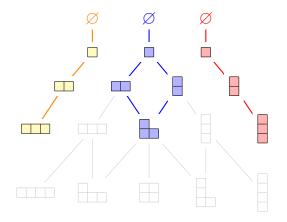


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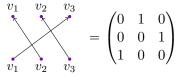
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Pick a basis for  $\mathbb{Q}^3$ :  $v_1 = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$   $v_2 = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}$   $v_3 = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$ Map each permutation to the matrix that permutes t

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& & & & \\
& & & & \\
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$$(13, 19, 22) = 13A + 19B + 22C$$

means A, B, and C got 13, 19, and 22 votes, respectively.

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Theorem. The permutation representation decomposes into one simple 1-dimensional (trivial) representation and one simple (n-1)-dimensional (reflection) representation.

The permutation representation models the outcome space. More: all voting data spaces are symmetric group "modules".

In other words:

Permutations naturally move around voting and outcome spaces.

(Permute the candidate's names, or their places on the ballot.)

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addition,

If individual precincts add up votes, and then combine results, that should be the same as if the tallying happened all in one place.

scaling, and

If everyone's vote counted 5 times, the outcome should be the same.

permutations.

Changing the order that candidates appear on the ballot ideally shouldn't change the outcome.

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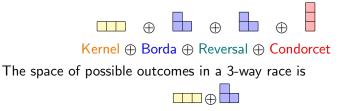
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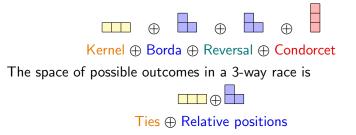
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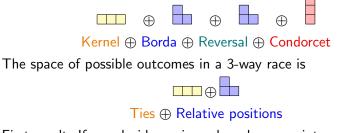
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Example: The space of possible votes in a 3-way race with full-rankings is



First result: If you decide a winner based on a points system, Condorcet cycles get lost in the tally.



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 $\mathsf{Kernel} \oplus \mathsf{Borda} \oplus \mathsf{Reversal} \oplus \mathsf{Condorcet}$ 

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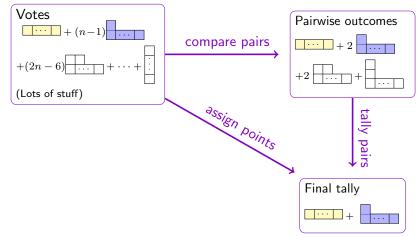
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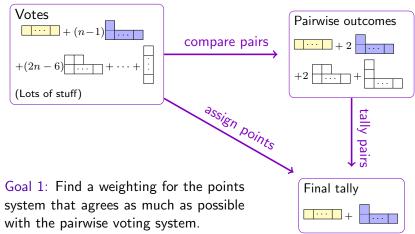
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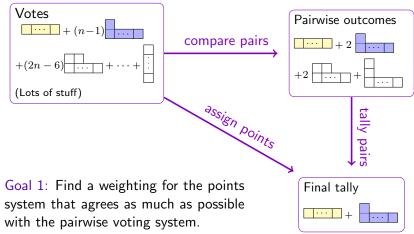


But by further analysis, one can compute that the image is at most

In fact, the information lost is precisely the "reversal" space.



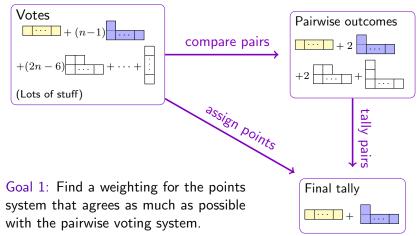




Result 1 (Saari, **D**.): The point-based system that agrees most with the pairwise outcomes is the (modified) Borda count, with weight

$$\mathbf{w} = \left(1, \ \frac{n-2}{n-1}, \ \dots, \ \frac{2}{n-1}, \ \frac{1}{n-1}, \ 0\right)$$

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Committees (Barcelo, Bernstein, Bockting-Conrad, McNichols, Nyman, Viel; '18): How to tally votes for committees with representation from several departments (using representation theory of  $S_m \wr S_n$ ).

## What do we want, and why do we care?

- Mathematicians: More data preservation is better.
- Idealists: Everyone should have a fair say.
- Pragmatists: Simpler voting systems are easier to implement.
- Cynics: Stupid things happen when people, en mass, are forced to game their votes.
- (Over-educated) Conspiracy Theorists: Our voting system is provably about as bad as it can be without *everyone* noticing.
- Kenneth Arrow: No voting system is ideal, so...
- Mathematicians (again): Oh, come on!

Some references...

• Many many publications of Donal Saari, particularly around 1999–2000. Also, "Decisions and Elections; Explaining the Unexpected", Cambridge University Press, 2001.

• "Voting, the symmetric group, and representation theory", by **D**., Eustis, Minton, and Orrison. American Mathematical Monthly 116 (2009), no. 8, 667–687.

• "Algebraic voting theory and representations of  $S_m \wr S_n$ ", by Barcelo, Bernstein, Bockting-Conrad, McNichols, Nyman, Viel. (preprint: arXiv:1807.03743)

"How not to be wrong", by Jordan Ellenberg.
 Chapter 17: "There is no such thing as public opinion."

And even though our system is non-ideal as is...

Go vote tomorrow! http://vote.nyc.ny.us