

An Algebraic Approach to Voting Theory

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Overview

- 1 Introduction to Voting Theory
 - Collecting Data
 - Tallying Data
 - Comparing Methods
- 2 Algebraic Methods
 - Tally Methods as Linear Maps
 - Relating Maps
- 3 Results
 - Full Rankings
 - Partial Rankings

Asking for votes

Full Rankings

Ex. 3 candidates

A	A	B	B	C	C
B	C	A	C	A	B
C	B	C	A	B	A

Partial Rankings

Ex. Top 3 from 5 candidates

A	A	A	A	A	A	A	A	A	A	A	A
B	B	B	C	C	C	D	D	D	E	E	E
C	C	E	B	D	E	B	C	E	B	C	D
DE	ED	CD	DE	BE	BD	CE	BE	BC	CD	BD	BC
⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮

Asking for votes

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⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮

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⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮

Tallying votes

Positional methods

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- Weighting vector

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$$\mathbf{w} = [1, 0, 0]$$

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$$\mathbf{w} = [2, 1, 0]$$

$$\mathbf{w} = [1, 1/2, 0]$$

- In general

$$\mathbf{w} = [1, t, 0], \quad 0 \leq t \leq 1$$

Tallying votes

Pairwise (Condorcet) method

A wins over B if $A > B$ more times than $B > A$.

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A wins over B if $A > B$ more times than $B > A$.

- Condorcet winner
- Condorcet criterion
- Cyclic preferences

Comparing methods

Example

<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>C</i>	<i>C</i>
<i>B</i>	<i>C</i>	<i>A</i>	<i>C</i>	<i>A</i>	<i>B</i>
<i>C</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
10	25	5	25	5	20

Comparing methods

Example

<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>C</i>	<i>C</i>
<i>B</i>	<i>C</i>	<i>A</i>	<i>C</i>	<i>A</i>	<i>B</i>
<i>C</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
10	25	5	25	5	20

Positional tally: $\mathbf{w} = [1, 0, 0]$ “Plurality”

Comparing methods

Example

<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>C</i>	<i>C</i>
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10	25	5	25	5	20

Positional tally: $\mathbf{w} = [1, 0, 0]$ “Plurality”

$$A : 10 + 25 = 35$$

$$B : 5 + 25 = 30$$

$$C : 5 + 20 = 25$$

Comparing methods

Example

<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>C</i>	<i>C</i>
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<i>C</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
10	25	5	25	5	20

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$$A : 10 + 25 = 35 \quad B : 5 + 25 = 30 \quad C : 5 + 20 = 25$$

Pairwise Tally

Comparing methods

Example

<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>C</i>	<i>C</i>
<i>B</i>	<i>C</i>	<i>A</i>	<i>C</i>	<i>A</i>	<i>B</i>
<i>C</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
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Pairwise Tally

$$B > A \quad 50 : 40$$

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<i>C</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
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Positional tally: $\mathbf{w} = [1, 0, 0]$ “Plurality”

$$A : 10 + 25 = 35 \quad B : 5 + 25 = 30 \quad C : 5 + 20 = 25$$

Pairwise Tally

$$B > A \quad 50 : 40 \quad C > A \quad 50 : 40 \quad C > B \quad 50 : 40$$

Pos: $A > B > C$ vs. Pairws: $C > B > A$

Comparing methods

Example

A	A	B	B	C	C
B	C	A	C	A	B
C	B	C	A	B	A
10	25	5	25	5	20

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<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>C</i>	<i>C</i>		
<i>B</i>	<i>C</i>	<i>A</i>	<i>C</i>	<i>A</i>	<i>B</i>		
<i>C</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>		
(10,	25,	5,	25,	5,	20)

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$$B > A \quad 50 : 40$$

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What is fair?

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 - **Non-dictatorship:** outcome dictated by more than one vote.
- There is no 'ideal' system.

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- Breaking down the profile space
- Borda Count minimizes conflicts

Tally Methods as Linear Maps

Positional Method

$$\mathbf{w} = [1, t, 0], \quad 0 \leq t \leq 1$$

$$T_{\mathbf{w}} = \begin{pmatrix} 1 & 1 & t & 0 & t & 0 \\ t & 0 & 1 & 1 & 0 & t \\ 0 & t & 0 & t & 1 & 1 \end{pmatrix} \begin{matrix} A \\ B \\ C \end{matrix}$$

Pairwise Method

$$P = \begin{pmatrix} 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 \end{pmatrix} \begin{matrix} A > B \\ A > C \\ B > A \\ B > C \\ C > A \\ C > B \end{matrix}$$

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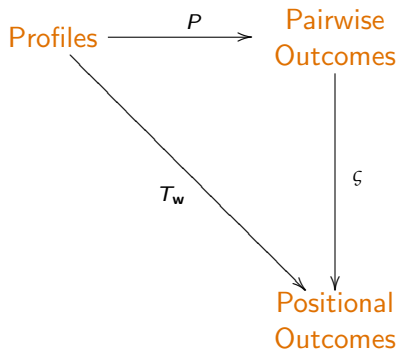
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Relating Tally Maps



Breaking down the spaces

Theorem

All pairwise and positional maps are $\mathbb{Q}S_n$ -module homomorphisms.

Breaking down the spaces

Theorem

All pairwise and positional maps are QS_n -module homomorphisms.

$$U \cong W_1 \oplus W_2 \oplus W_3 \oplus Y_1 \oplus Z_1$$

$$V \cong W_4 \oplus W_5 \oplus Y_2 \oplus Y_3.$$

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Theorem

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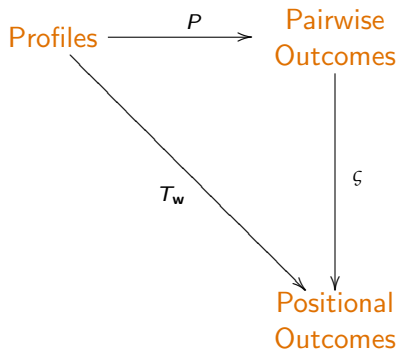
$$W_2 \mapsto W_5$$

$$W_3 \mapsto 0$$

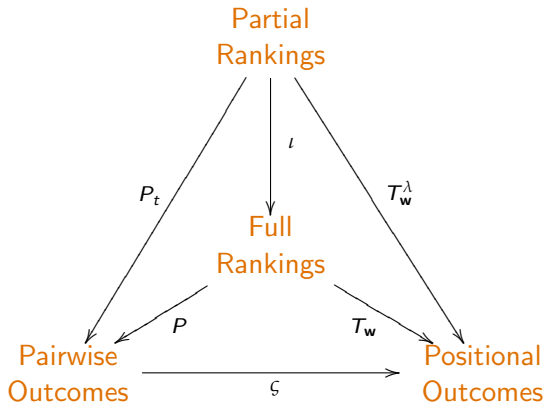
$$Y_1 \mapsto Y_2$$

$$Z_1 \mapsto 0$$

Relating Tally Maps



More Comparisons



Results

Full Rankings:

Results

Full Rankings:

- Recovery of Saari's results

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- Borda Count uniquely allows our maps to commute

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Full Rankings:

- Recovery of Saari's results
- Borda Count uniquely allows our maps to commute
- Tools for a new perspective

Results

For rankings of top k from n candidates, two cases:

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- If pairwise map gives tying candidates each $1/2$ point, a particular scaled linear modification of the Borda Count works uniquely.

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$$\begin{array}{c} A \\ B \\ C \\ \hline D \\ E \end{array}$$

Results

For rankings of top k from n candidates, two cases:

- If pairwise map gives tying candidates each $1/2$ point, a particular scaled linear modification of the Borda Count works uniquely.

$$\begin{array}{r} A \quad 1 \\ B \quad 3/4 \\ C \quad 2/4 \\ \hline D \quad 1/4 \\ E \quad 0 \end{array}$$

Results

For rankings of top k from n candidates, two cases:

- If pairwise map gives tying candidates each $1/2$ point, a particular scaled linear modification of the Borda Count works uniquely.

A	1	1
B	$3/4$	$3/4$
C	$2/4$	$2/4$
<hr/>		
D	$1/4$	$1/8$
E	0	$1/8$

Results

For rankings of top k from n candidates, two cases:

- If pairwise map gives tying candidates each $1/2$ point, a particular scaled linear modification of the Borda Count works uniquely.

A	1	1	$7/8$
B	$3/4$	$3/4$	$5/8$
C	$2/4$	$2/4$	$3/8$
<hr/>			
D	$1/4$	$1/8$	0
E	0	$1/8$	0

Results

For rankings of top k from n candidates, two cases:

- If pairwise map gives tying candidates each $1/2$ point, a particular scaled linear modification of the Borda Count works uniquely.

A	1	1	$7/8$	1
B	$3/4$	$3/4$	$5/8$	$5/7$
C	$2/4$	$2/4$	$3/8$	$3/7$
D	$1/4$	$1/8$	0	0
E	0	$1/8$	0	0

Results

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- If pairwise map gives tying candidates each $1/2$ point, a particular scaled linear modification of the Borda Count works uniquely.

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C	$2/4$	$2/4$	$3/8$	$3/7$
<hr/>				
D	$1/4$	$1/8$	0	0
E	0	$1/8$	0	0

- If pairwise map gives tying candidates each something besides $1/2$ point, there may be more freedom in our choice of positional method.

Thanks to...

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Professor Michael Orrison

Fellow thesis students

for more information:

<http://www.math.hmc.edu/~zajj/thesis/>