MATH 113: DISCRETE STRUCTURES WEDNESDAY WEEK 11 HANDOUT

The greatest common divisor d = gcd(a, b) of integers a, b is the largest positive integer such that $d \mid a$ and $d \mid b$. We say that a and b are *relatively prime* when they share no divisors larger than 1, and this is equivalent to gcd(a, b) = 1.

Problem 1. Draw a divisor diagram for 84 and 105. Where does the gcd appear in partially ordered set of divisors?

If we know the prime factorizations of *a* and *b*, this number is easy to determine. Let $\{p_1, p_2, ..., p_k\}$ be the set of distinct prime divisors of *a* and *b*. Then we may write

$$a = p_1^{a_1} p_2^{a_2} \cdots p_k^{a_k},$$

$$b = p_1^{b_1} p_2^{b_2} \cdots p_k^{a_k}$$

for nonnegative integers a_i , b_i and

$$gcd(a,b) = p_1^{\min\{a_1,b_1\}} p_2^{\min\{a_2,b_2\}} \cdots p_k^{\min\{a_k,b_k\}}$$

It is frequently the case, though, that we do not have access to the prime factorizations of integers. In this case, the *Euclidean algorithm* allows us to determine the greatest common divisor. Let's execute the algorithm with a = 81, b = 57:

$$81 = 1 \cdot 57 + 24$$

$$57 = 2 \cdot 24 + 9$$

$$24 = 2 \cdot 9 + 6$$

$$9 = 1 \cdot 6 + 3$$

$$6 = 2 \cdot 3 + 0.$$

We conclude that the final nonzero remainder, 3, is the gcd of 81 and 57. Indeed, $81 = 3^4$ and $57 = 3 \cdot 19$, so this agrees with our first method for determining gcd's.

The Euclidean algorithm can be described formally as follows:

- 1. Assume a > b are integers (if a < b, swap them).
- 2. Perform long division to express to express a = qb + r where $0 \le r \le b 1$.
- 3. Replace a with b and b with r.
- 4. If $r \neq 0$, return to step 2; else
- 5. if r = 0, conclude that the final nonzero remainder is gcd(a, b).

A generic run of the algorithm then looks like

$$a = q_0 b + r_1$$

$$b = q_1 r_1 + r_2$$

$$r_1 = q_2 r_2 + r_3$$

$$r_2 = q_3 r_3 + r_4$$

:

$$r_{n-2} = q_{n-1} r_{n-1} + r_n$$

$$r_{n-1} = q_n r_n + 0$$

where $1 \le r_k \le r_{k-1}$ and we conclude that $r_n = \text{gcd}(a, b)$ (since $r_{n+1} = 0$).

Problem 2. Suppose an integer *x* divides integers *y* and *z*. Show that for any $k, \ell \in \mathbb{Z}$, $x \mid ky + \ell z$.

Problem 3. Why does the Euclidean algorithm work? Start at the end of the algorithm and check that $r_n | r_{n-1}$, then inductively check that $r_k | r_{k-1}$ for $-1 \le k \le n$ where we write $r_0 = b$ and $r_{-1} = a$ for notational convenience. Conclude that r_n divides a and b. Use a similar argument starting at the beginning of the algorithm to show that gcd(a, b) divides r_k for $-1 \le k \le n$. Why does this prove that the algorithm produces the gcd.

Problem 4. The Euclidean algorithm gives us a way to dissect a rectangle with integer sides into squares. Run the Euclidean algorithm to find gcd(23, 13). Interpret the first step $(23 = 1 \cdot 13 + 10)$ as telling you that $q_0 = 1$ -many 10×10 squares fit inside a 23×13 rectangle. Figure out what instructions the rest of the algorithm is giving you and draw a corresponding picture. At the end, your 23×13 rectangle should be partitioned into squares! What is special about this procedure if you start with consecutive Fibonacci numbers $a = F_{n+1}$, $b = F_n$?

Problem 5. Run the Euclidean algorithm when a = 45, b = 16. How is it related to the expression

$$\frac{45}{16} = 2 + \frac{1}{1 + \frac{1}{4 + \frac{1}{3}}}?$$

Come up with a general procedure by which the Euclidean algorithm produces *continued fraction* expressions for rational numbers of the form

$$\frac{a}{b} = x_1 + \frac{1}{x_2 + \frac{1}{x_3 + \frac{1}{x_4 + \dots}}}$$

where the x_i are integers.