PROBLEM 1. For each of the following, prove either that the language is regular or that it is not regular. In all cases $\Sigma = \{0.1\}$.

- (a) $L_1 = \{w : w \text{ contains an equal number of 0s and 1s}\}.$
- (b) $L_2 = \{1^k w : w \in \Sigma^*, k \ge 1, \text{ the number of 1s in } w \text{ is at least } k\}.$
- (c) $L_3 = \{1^k w : w \in \Sigma^*, k \ge 1, \text{ the number of 1s in } w \text{ is at most } k\}.$

PROBLEM 2. Consider the context-free grammar G

$$\begin{split} S &\to RR \\ R &\to 0R1 \mid \varepsilon. \end{split}$$

- (a) Which of the following are in L(G)? For each that is, give a left-most derivation making one substitution at each step.
 - (i) 001101
 - (ii) 000111.
 - (iii) 1010.
 - (iv) ε .
- (b) Succinctly describe the language L(G). Explain your reasoning.
- (c) Is L(G) regular? Provide a proof.

Solution.

$$\begin{array}{ll} (\mathrm{a}) & (\mathrm{i}) \ 001101 \in L(G): \\ & S \rightarrow RR \rightarrow 0R1R \rightarrow 00R11R \rightarrow 0011R \rightarrow 00110R1 \rightarrow 001101. \\ (\mathrm{ii}) \ 000111 \in L(G): \\ & S \rightarrow RR \rightarrow R \rightarrow 0R1 \rightarrow 00R11 \rightarrow 000R111. \\ (\mathrm{iii}) \ 1010 \not\in L(G). \\ (\mathrm{iv}) \ \varepsilon \in L(G): \\ & S \rightarrow RR \rightarrow R \rightarrow \varepsilon. \end{array}$$

(b) The language produced by the CFG

$$S \to R$$
$$R \to 0R1|\varepsilon$$

is $A := \{0^k 1^k : k \ge 0\}$. The production $S \to RR$ in G then says that $L(G) = A^2 = AA := \{ww : w \in A\}$.

(c) The language L(G) is not regular. We prove this by contradiction. If it we regular, then take a pumping length p and consider the word $0^p 1^p \in L(G)$. By the pumping lemma $0^p 1^p = xyz$ where $|xy| \leq p$, |y| > 0, and $xy^i z \in L(G)$ for all $i \geq 0$. Then yconsists of all 0s, and pumping it would produce words in L(G) that have more 0s than 1s, which is not possible. PROBLEM 3. Let A be the language of all words in $\Sigma = \{0, 1\}$ that contain the same number of occurrences of 01 as occurrences of 10. Is A regular? Provide a proof.

Solution. The language A is regular. In fact, A is given by the regular expression $(0^+1^+0^+)^* \cup (1^+0^+1^+)^*$.

PROBLEM 4. Let A be the NFA pictured below:



- (a) Create a CFG G such that L(A) = L(G).
- (b) Give a derivation of the word *aabab* making one substitution at each step.

PROBLEM 5. Let $P = \{0^q : q \text{ is a prime number}\}$. Prove or disprove that P is regular.

PROBLEM 6. (Bonus) Our goal in this problem is to show that the representation of objects can affect whether or not a given set can be recognized by a machine. Consider a set Aof natural numbers. Let $B_k(A)$ be the set of strings that represent numbers from A in base k (with no leading zeros). For example, if $A = \{3, 5\}$ then $B_2(A) = \{11, 101\}$ and $B_3(A) = \{10, 12\}$. We can think of $B_k(A)$ as a language with a k-symbol alphabet. Give a set A for which $B_2(A)$ is regular but $B_3(A)$ is not (and prove it).