Homework 8F 4(a)

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Proposition (Problem 4(a)'s tricky direction). Let λ be an element of k^{\times} , and let q be a regular quadratic form over k with dim q = 2m. Then, $q \cong \langle \lambda \rangle \otimes q \implies q \cong q_1 \perp \cdots \perp q_m$, where each q_i is a binary form such that $q_i \cong \langle \lambda \rangle \otimes q_i$.

Proof. We will show this by induction on the dimension (with an induction step size of 2). For the base case when dim q = 2, the equality is trivially established.

Suppose that for dim q' = 2(m-1), $q' \cong \langle \lambda \rangle \otimes q' \implies q' \cong q_1 \perp \ldots \perp q_{m-1}$ in which $q_i \cong \langle \lambda \rangle \otimes q_i$ are binary forms $\forall i < m-1$. Let's find the case when dim q = 2m. Pick some diagonalization so that $q = \langle \mu_1, \ldots, \mu_{2m} \rangle$ and $\langle \lambda \rangle \otimes q = \langle \lambda \mu_1, \ldots, \lambda \mu_{2m} \rangle$. If two quadratic forms $q \cong \langle \lambda \rangle \otimes q$, then, by Witt's chain equivalence theorem, \exists a binary subform¹ f of q and a binary subform g of $\langle \lambda \rangle \otimes q$ such that $f \cong g$. This follows from the existence of chain equivalence and the definition of simple equivalence.² This means that $\exists i, j, k, l \in \mathbb{N}$ such that $\langle \mu_i, \mu_j \rangle \cong \langle \lambda \mu_l, \lambda \mu_k \rangle$. However, we are allowed to permute the diagonal elements to different positions while leaving the form isometric. Hence, WLOG, we can demand that i = l = 2m - 1 and j = k = 2m, *i.e.* $\langle \lambda \rangle \otimes q_m \cong q_m$ in which $q_m = \langle \mu_{2m-1}, \mu_{2m} \rangle$. Then, by the Witt's cancellation theorem, we can cancel out this part and find that $\langle \mu_1, \ldots, \mu_{2(m-1)} \rangle \cong \langle \lambda \mu_1, \ldots, \lambda \mu_{2(m-1)} \rangle$. Invoking the induction hypothesis, we can infer that the proposition holds for the dim q = 2m case.

Remark. What I learned from this proof is that induction proofs for quadratic forms' equivalence have a "natural step size" of 2 due to the combination of the chain equivalence theorem and the cancellation theorem.

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¹ By a binary subform of q, we mean, quite intuitively, a binary quadratic form that is a part of the orthogonal sum of q.

² Note that simple equivalence also permits equivalence between unary subforms, but it that case we can always group two unary subforms into a binary subform.