## MATH 113: DISCRETE STRUCTURES MONDAY WEEK 4 SUPPLEMENT

Let's first show that the identity

$$1 \cdot n + 2 \cdot (n-1) + 3 \cdot (n-2) + \dots + (n-1) \cdot 2 + n \cdot 1 = \binom{n+2}{3}$$

holds for integers  $n \ge 2$ .

*Proof 1 (double counting).* Observe that the right-hand side gives the number of 3-subsets of  $\underline{n+2} = \{1, 2, ..., n+2\}$ . We show that the sum on the left-hand side counts the same objects.

Given a 3-subset *A* of  $\underline{n+2}$ , write it as  $\{a, b, c\}$  where a < b < c by ordering its elements. For a given integer *b* between 2 and n + 1, we count the number of 3-subsets with *b* as their middle element. The first element, *a*, can be any of the b - 1 elements of  $\{1, 2, ..., b - 1\}$ , while the final element, *c*, can be any of the n+2-b elements of  $\{b+1, b+2, ..., n+2\}$ . (Stop and think for a while if it is not immediately clear that there are n+2-b such elements.) Thus there are (b-1)(n+2-b)3-subsets of  $\underline{n+2}$  with *b* as middle element. Letting *b* range from 2 (the smallest possible middle element) to n + 1 (the largest possible middle element), we count a total of

$$1 \cdot n + 2 \cdot (n-1) + \dots + (n-1) \cdot 2 + n \cdot 1$$

3-subsets of n + 2. (Note that we have not double-counted anything because subsets with distinct middle elements are distinct.) We conclude that

$$1 \cdot n + 2 \cdot (n-1) + 3 \cdot (n-2) + \dots + (n-1) \cdot 2 + n \cdot 1 = \binom{n+2}{3}$$

as desired.

*Proof 2 (induction).* We first check the base case, n = 2. For this n, the left-hand side is  $1 \cdot 2 + 2 \cdot 1 = 4$  and the right-hand side is  $\binom{4}{3} = 4$ , so the identity holds.

Now assume that for some  $n \ge 2$  we have

$$1 \cdot n + 2 \cdot (n-1) + 3 \cdot (n-2) + \dots + (n-1) \cdot 2 + n \cdot 1 = \binom{n+2}{3}.$$

Replacing *n* with n + 1 on the left-hand side we get

$$1 \cdot (n+1) + 2 \cdot (n+1-1) + 3 \cdot (n+1-2) + \dots + (n+1-1) \cdot 2 + (n+1) \cdot 1$$

Here each term is of the form  $k \cdot (n + 2 - k)$  where k ranges from 1 to n + 1. Note that we can rewrite  $k \cdot (n + 2 - k)$  as  $k \cdot (1 + (n + 1 - k)) = k + k \cdot (n + 1 - k)$ . Adding these terms up (first the k summand, then the  $k \cdot (n + 1 - k)$  summand), we get

$$(1+2+\dots+n+(n+1))+(1\cdot n+2\cdot (n-1)+3\cdot (n-2)+\dots+(n-1)\cdot 2+n\cdot 1).$$

By the induction hypothesis, the second term is just  $\binom{n+2}{3}$ . By the combinatorial argument we gave in class, the first term is  $\binom{n+2}{2}$ . We can thus conclude that

$$1 \cdot (n+1) + 2 \cdot (n+1-1) + 3 \cdot (n+1-2) + \dots + (n+1-1) \cdot 2 + (n+1) \cdot 1 = \binom{n+2}{2} + \binom{n+2}{3}.$$

Finally, by Pascal's identity, the final sum is  $\binom{n+3}{3} = \binom{(n+1)+2}{3}$ , as desired. By mathematical induction, we now know that the identity holds for all  $n \ge 2$ .

Let's now show that for  $1 \le k \le n$ ,

$$n^k \le \binom{n}{k} k^k.$$

Before getting into the proof, note that if we divide both sides by  $k^k$  (a positive number), we get

$$\frac{n^k}{k^k} \le \binom{n}{k},$$

which is equivalent to

$$\left(\frac{n}{k}\right)^k \le \binom{n}{k}$$

for  $1 \le k \le n$ , which, if nothing else, is an orthographically attractive inequality. (*Challenge*: Run some computer experiments to see how effective the bound  $(n/k)^k$  is on  $\binom{n}{k}$ . Are they close or far apart as n gets large?)

*Proof.* Our strategy is as follows: given  $1 \le k \le n$ , we construct sets X and Y such that  $|X| = \binom{n}{k}k^k$  and  $|Y| = n^k$ . We will then exhibit a surjection  $f : X \to Y$ . This guarantees the inequality  $|X| \ge |Y|$ , which is equivalent to  $n^k \le \binom{n}{k}k^k$ .

Fix integers k and n with  $1 \le k \le n$ . We take Y to be the collection of words in  $\underline{n} = \{1, ..., n\}$  of length k. In other words,

$$Y = \{(a_1, a_2, \dots, a_k) \mid a_1, \dots, a_k \in \underline{n}\}.$$

Since there are *n* choices for each letter in each length *k* word, we have  $|Y| = n^k$ .

Now let *X* be the collection of pairs  $(A, (a_1, \ldots, a_k))$  where *A* is a *k*-subset of  $\underline{n}$  and  $a_1, \ldots, a_k \in A$ . There are  $\binom{n}{k}$  ways to select the first term, *A*, and then  $k^k$  possible length *k* words  $(a_1, \ldots, a_k)$  constructed from elements of *A*. Thus  $|X| = \binom{n}{k}k^k$ .

Finally, we construct a surjective function  $f: X \to Y$  by declaring that

$$f((A, (a_1, \ldots, a_k))) = (a_1, \ldots, a_k).$$

This is well-defined for any  $(A, (a_1, \ldots, a_k)) \in X$  as the word  $(a_1, \ldots, a_k)$  is constructed from a subset of  $\underline{n}$ . Moreover, for each word  $(a_1, \ldots, a_k) \in Y$ , we can construct a pair  $(A, (a_1, \ldots, a_k)) \in X$  such that  $f((A, (a_1, \ldots, a_k))) = (a_1, \ldots, a_k)$  by choosing some A which contains  $\{a_1, \ldots, a_k\}$ . This means that f is surjective, so we have finished our proof!