Algebraic geometry of sandpiles

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# PCMI 2008 Undergraduate Summer School Lecture 15: Sandpiles

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- Self-Organized Criticality: An Explanation of 1/f Noise, Bak, Tang, Wiesenfeld, Physical Review Letters, 1987.
- Self-Organized Critical State of Sandpile Automaton Models, Dhar, Physical Review Letters, 1990.

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# Abelian Sandpile Model

#### Notation

G = (V, E, s): finite, connected, loopless, multigraph with a distinguished vertex *s* called the sink

sandpile configurations:  $\mathbb{N}^{V_s}$  where  $V_s = V \setminus \{s\}$ 

monoid structure:  $(c+c')_v = c_v + c'_v$  for all  $v \in V_s$ 

The Sandpile Group	Algebraic geometry of sandpiles	Tilings 00000	Complexity o
Stability			

$$\deg(v) = |\{\{v, w\} \in E : w \in V\}|$$

 $v \in V_s$  is stable in a configuration c if  $c_v < \deg(v)$ 

## Toppling

If v is an unstable vertex in c, we get a new configuration c' by toppling c at v:

$$c'_w = \left\{ egin{array}{cl} c_w + 1 & ext{if } \{v,w\} \in E \ c_w - \deg(w) & ext{if } w = v \ c_w & ext{otherwise} \end{array} 
ight.$$

- By a series of topplings, every configuration reaches a stable configuration.
- This stable configuration is independent of the ordering of the topplings.

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# Monoid of Stable Sandpiles

# Let $\ensuremath{\mathcal{S}}$ denote the commutative monoid of stable sandpile configurations with addition define by

$$c \circledast c' = \text{stabilization}(c + c').$$

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# **Recurrent Configurations**

## Definition

A stable configuration c is recurrent if given any configuration c', there exists a configuration c'' such that

stabilization(c' + c'') = c.

#### Example

Define  $c_{\max}$  by  $(c_{\max})_{\nu} = \deg \nu - 1$ . Then  $c_{\max}$  is recurrent.

#### HW

The recurrent configurations are exactly configurations that can be reached by adding sand to  $c_{\rm max}$  and stabilizing.

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# The Sandpile Group

#### Theorem/Definition

The collection of recurrent configurations,  $\mathcal{G}$  is a group called the sandpile group.

Interesting question: What is the identity of *G*? Algebraic geometry of sandpiles

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# Laplacian

Let 
$$V_s = \{v_1, ..., v_n\}.$$

#### Definition

The reduced Laplacian matrix for G is the  $n \times n$  matrix, L, where

$$\mathcal{L}_{ij} = \left\{ egin{array}{ll} \deg v_i & ext{if } i=j \ -m & ext{if } \{i,j\} \in E, \ m ext{ times} \ 0 & ext{otherwise} \end{array} 
ight.$$

#### Note

If c is a configuration with unstable vertex  $v_i$ , then

$$c - Le_i$$

is the configuration obtained by toppling  $v_i$ .

The	Sand	pile	Group
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#### Theorem

## $\mathcal{G} \approx \mathbb{Z}^n / \textit{image}(L)$

#### Corollary

$$|\mathcal{G}| = \det L = the number of spanning trees of G$$

Proof. Matrix-tree theorem.

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Let 
$$V = \{v_1, ..., v_{n+1}\}$$
 where

- $v_{n+1} = s$ , the sink.
- If  $v_i$  is farther than  $v_j$  from the sink, then i > j.

Label  $v_i$  with the indeterminate  $x_i$  for each *i*.

For 
$$i \in \{1, ..., n + 1\}$$
,  
 $p_i = x_i^{\deg v_i} - \prod_{j: \{v_i, v_j\} \in E} x_j$ , setting  $x_{n+1} = 1$ .

#### Definition

The sandpile ideal for G is

$$I_G = (p_i : i = 1, \dots, n+1) \subseteq \mathbb{C}[x_1, \dots, x_n]$$

Consider  $\mathbb{C}[x_1, \ldots, x_n]$  with graded revlex monomial ordering and

 $x_1 > \cdots > x_n$ .

## Theorem (Cori, Rossin, Salvy, 2006)

A normal basis for  $\mathbb{C}[x_1, \ldots, x_n]/I_G$  with respect to the above ordering is in one-to-one correspondence with the elements of the sandpile group,  $\mathcal{G}$ :

$$x^e \leftrightarrow C_{\max} - e.$$

#### Theorem (P., 2008)

 $Z(I_G) \subset \mathbb{A}^n$  is a set of  $|\mathcal{G}|$  points forming an orbit of a representation of  $\mathcal{G}$  by a group of  $n \times n$  matrices.

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## Theorem (Baker, et al., 2007)

There is more than an analogy between algebraic curves and graphs: Riemann-Roch, Riemann-Hurwitz, Jacobi inversion, etc. The sandpile group plays the role of the Picard group.

#### Corollary

The postulation number for the (homgenization of the ) sandpile ideal is the genus of the graph: g = |E| - |V| + 1.

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Order of the all-2s sandpile on an even square grid Let  $g_n$  be the order of the all-2s element of a  $2n \times 2n$  grid. The first few values of  $g_n$  are

*g*<sub>n</sub> : 1, 3, 29, 901, 89893

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#### Tilings of the even square grid

Let  $T_n$  be the number of domino tilings of a  $2n \times 2n$  grid.

$$T_n = 4^{n^2} \prod_{i,j=1}^n \left( \cos^2 \frac{i\pi}{2n+1} + \cos^2 \frac{j\pi}{2n+1} \right)$$
  
=  $2^n a_n^2$ 

where the first few values of  $a_n$  are

 $a_n$ : 1, 3, 29, 901, 89893

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#### A few more values of $a_n$ and $g_n$

an	1	3	29	901	89,893	28,793,575	29,607,089,625
g <sub>n</sub>	1	3	29	901	89,893	5,758,715	22,687,425

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## Theorem (Morar, P., 2007)

For each n, there are graphs  $P_n$  and  $A_n$  such that

- The number of domino tilings (perfect matchings) of A<sub>n</sub> is a<sub>n</sub>.
- The spanning trees of *P<sub>n</sub>* are in 1-1 correspondence with the domino tilings of *A<sub>n</sub>*.
- The cyclic subgroup generated by the all-2s element injects into the sandpile group of P<sub>n</sub>.

Thus,  $g_n$  divides  $a_n$ .

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## Theorem (Morar, P., 2007)

Let  $F_i$  be the *i*-th Fibonnaci number (starting with  $F_1 = F_2 = 1$ ). The order of the all-2s element for the 2 × n grid is:

for n = 2m,

order(*all 2s*) = 
$$\begin{cases} F_{n+1}/2 & \text{if } 3|(m-1)\\ F_{n+1} & \text{otherwise.} \end{cases}$$

2 for n = 2m - 1,

order(*all-2s*) = 
$$\begin{cases} (F_n + F_{n+2})/2 & \text{if } 3|m \\ F_n + F_{n+2} & \text{otherwise.} \end{cases}$$

In particular, the order of the all-2s element is odd.

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*Sandpile as a universal computer*, Goles and Margenstern, International Journal of Modern Physics, 1996.

*Constructing a sandpile machine*, Schoenberg-Jones, Reed College Thesis, 2008.